

Part 4. Queues

CS 200 Algorithms and Data Structures

Outline

- **Queue?**
- Implementing Queue
- Comparison implementations

2



Photo by David Jump

"Grill the Buffs" event 9/16 2011



4

Queue

- A *queue* is like a *line of people*.
- New item enters a queue at its back.
- Items leave a queue from its front.
- **First-in, first-out (FIFO)** behavior
- Removing and adding are done from *opposite ends* of structure
- Useful for scheduling (e.g. print queue, job queue)



5

Operations

- **Create** an empty queue
- Determine whether a queue is **empty**
- **Add** a new item to the queue
- **Remove** item from the queue (that was added the *earliest*)
- **Remove all** items from the queue
- **Retrieve** item from queue that was added earliest

6

Queue Operations

- **enqueue**(in newItem: QueueItemType)
 - Add new item at the back of a queue
- **dequeue**() ; QueueItemType
 - Retrieves and removes the item at the *front* of a queue
- **peek**() : QueueItemType {query}
 - Retrieve item from the *front* of the queue. Retrieve the item that was added earliest.
- **isEmpty**() : boolean {query}
- **createQueue**()
- **dequeueAll**()

7

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- **Implementing Queue**
- Comparison implementations

8

Implementations of Queue

- *Reference-Based* Implementation
- *Array-based* Implementation
- *List based* Implementation
- *java.util.queue* interface

9

Reference-based Implementation(1)

- Needs
 - **Nodes** with the *item* and a *reference* to the next item
 - **two external references** pointing to the *first* node and the *last* node.

10

Reference-based Implementation(2)

- **Single external references**
 - *Circular linked list* represents a queue
 - The node at the back of the queue references the node at the front

11

Inserting an item into a nonempty queue

- Step 1. `newNode.next = lastNode.next;`
- Step 2. `lastNode.next = newNode;`
- Step 3. `lastNode = newNode;`

12

isEmpty()

```
public class QueuerReferenceBased implements
QueueInterface {
    private Node lastNode;
    public QueuerReferenceBased(){
        lastNode = null;
    }
    public boolean isEmpty(){
        return lastNode == null;
    }
}
```

13

Insert new item into the queue

```
public void enqueue (Object newItem){
    Node newNode = new Node(newItem);
    if (isEmpty()){
        newNode.next = newNode;
    } else {
        newNode.next = lastNode.next;
        lastNode.next = newNode;
    }
    lastNode = newNode;
}
```

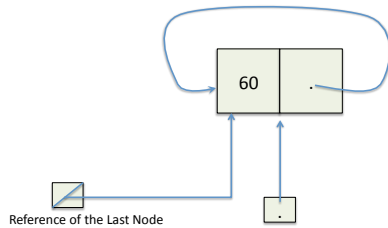
A. Empty queue

B. More than 1 item

14

Inserting a New Item

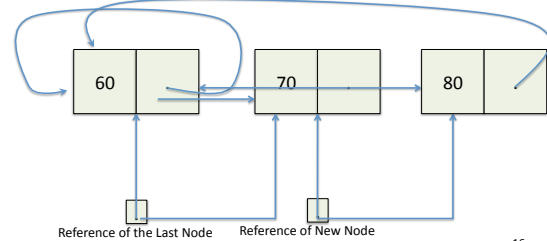
- Insert a **new item** into the **empty queue**



15

Inserting a New Item

- Insert a **new item** into a queue that has **more than 1 items**.



16

Removing an item from queue

```
public Object dequeue() throws QueueException{
    if (!isEmpty()){
        Node firstNode = lastNode.next;
        if (firstNode == lastNode) {
            lastNode = null;
        }
        else{
            lastNode.next = firstNode.next;
        }
        return firstNode.item;
    }
    else { exception handling..
    }
}
```

Why?

17

Peek?

```
public Object peek() throws QueueException {
    if (!isEmpty()){
        Node firstNode = lastNode.next;
        return firstNode.item;
    }else {
        throw new QueueException(your_message);
    }
}
```

18

Implementing queue with Array

- We need,
 - An **array** of objects to store items
 - variables** to point to the "front" and "back" index of the array

The diagram shows two boxes labeled 'Front' and 'Back' containing the values 0 and 1 respectively. To their right is an array labeled 'items' with indices 0, 1, 2, 3, and MAX_QUEUE-1. The array contains the values 2 and 4 at indices 0 and 1, and is empty at index MAX_QUEUE-1.

Implementing queue with Array

- Queue is **empty** if *back* is less than *front*
- Inserting an item
 - Initially *front* is 0, *back* is -1
 - Increment *back*
 - Place new item in *items[back]*
- Queue will be **full** if *back* equals $\text{MAX_QUEUE} - 1$

Implementing queue with Array

- Removing** an item
 - Remove item from *items[front]*
 - Increment front**
- Rightward drift**
 - After a sequence of additions and removals, items in the queue will **drift toward the end** of the array
 - back* can reach $\text{MAX_QUEUE} - 1$ even when the queue contains only few items.

Solving Rightward Drift Problem (1)

- Shifting** array elements to **the left**. Therefore front is always pointing to *items[0]*.
- Cost** of implementation

Solving Rightward Drift Problem (2)

- Circular implementation** of a queue

The diagram shows a circular array with 7 slots labeled 0 through 6. An arrow labeled 'FRONT' points to slot 0, and an arrow labeled 'BACK' points to slot 3. Slots 0, 1, 2, and 3 contain items 2, 4, 1, and 7 respectively.

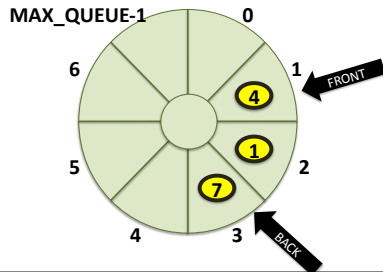
Solving Rightward Drift Problem (2)

- Delete**

The diagram is identical to the previous one, showing a circular queue with items 2, 4, 1, 7 at indices 0, 1, 2, 3. The 'FRONT' arrow is at index 0 and the 'BACK' arrow is at index 3.

Solving Rightward Drift Problem (2)

• Delete

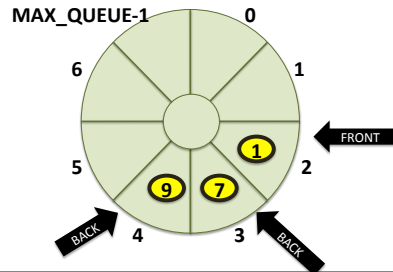


25

Solving Rightward Drift Problem (2)

• Insert 9

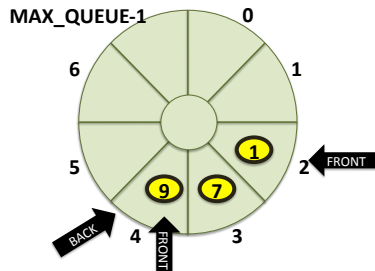
When either front or back advances past MAX_QUEUE-1, it wraps around 0



26

Queue with Single Item

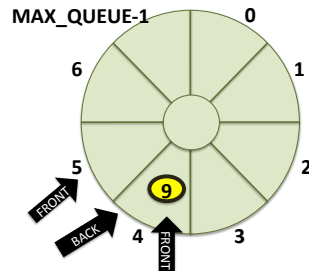
- Queue contains **only one item**.
- *back* and *front* are pointing at the same slot.



27

Queue with Single Item

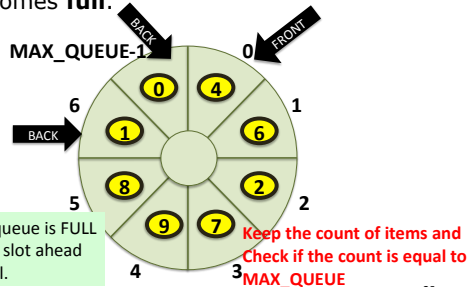
- The last item is removed. Queue is empty. - *front* passed *back*.



28

Insert the last item

- *back* catches up to *front* when the queue becomes **full**.



When the queue is FULL *front* is one slot ahead *back* as well.

Keep the count of items and Check if the count is equal to MAX_QUEUE

29

Wrapping the values for front and back

- Initializing
 - `Front = 0`
 - `Back = MAX_QUEUE-1`
 - `Count = 0`
- Adding


```
back = (back+1) % MAX_QUEUE;
items[back] = newItem;
++count;
```
- Deleting


```
deleteItem = items[front];
front = (front + 1) % MAX_QUEUE;
--count;
```

30

enqueue with Array

```
public void enqueue(Object newItem) throws
QueueException{
    if (!isFull()){
        back = (back+1) % (MAX_QUEUE);
        items[back] = newItem;
        ++count;
    }else {
        throw QueueException(your_message);
    }
}
```

31

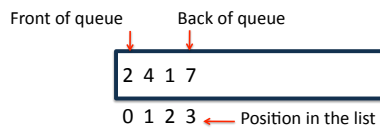
dequeue()

```
public Object dequeue() throws QueueException{
    if (!isEmpty()){
        Object queueFront = items[front];
        front = (front+1) % (MAX_QUEUE);
        --count;
        return queueFront;
    }else{
        throw new QueueException (your_message);
    }
}
```

32

Implementation with List

- You can implement operation **dequeue()** as the list operation **remove(0)**.
- **peek()** as **get(0)**
- **enqueue()** as **add (size()-1, newItem)**



33

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34

java.util.Queue

- extends `java.util.Collection`
- `add(e)`, `remove()`, `element()`
- `offer(e)`, `poll()`, `peek()`
- `java.util.Deque`:
 - Subinterface of queue
 - A linear collection that supports **element insertion and removal at both ends**.

35

Summary of Queue Operations (FIFO)

ADT Queue	JCF Queue	JCF Deque
enqueue(e)	add(e)	addLast(e)
	offer(e)	offerLast(e)
dequeue()	remove()	removeFirst()
	poll()	pollFirst()
peek()	peek()	peekFirst()
	element()	getFirst()

36

Summary of LIFO operations

ADT Stack	JCF Stack	JCF Deque
push(e)	push(e)	addFirst(e)
pop()	pop()	removeFirst()
peek()	peek()	peekFirst()

37

Summary of Position-Oriented ADTs

- **Stack**, **queue**, and **list** (so far)
- createStack and createQueue
- isEmpty for stack and queue
- push and enqueue
- pop and dequeue
- Stack peek and queue peek

38

Next Reading

- Section 9. Advanced Java Topics

39