

# Reference Sheet for Logic and Program Proofs

## Logical Equivalences

Definition of $\wedge$ $P \wedge \neg P \equiv \text{False}$ $P \wedge \text{False} \equiv \text{False}$ $P \wedge \text{True} \equiv P$	Idempotent Laws $p \vee p \equiv p$ $p \wedge p \equiv p$	DeMorgan's Laws $\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$	Distributive Laws $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$
Definition of $\vee$ $P \vee \neg P \equiv \text{True}$ $P \vee \text{False} \equiv P$ $P \vee \text{True} \equiv \text{True}$	Double Negation $\neg(\neg p) \equiv p$	Absorption Laws $p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Associative Laws $(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
	Commutative Laws $p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Implication Laws $p \rightarrow q \equiv \neg p \vee q$ $p \rightarrow q \equiv \neg q \rightarrow \neg p$	Biconditional Laws $p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$ $p \leftrightarrow q \equiv \neg q \leftrightarrow \neg p$

## Inference Rules

Simplification $p \wedge q$ <hr/> Therefore, $p$	Modus Ponens $p$ $p \rightarrow q$ <hr/> Therefore, $q$	Modus Tollens $\neg q$ $p \rightarrow q$ <hr/> Therefore, $\neg p$	Hypothetical Syllogism $p \rightarrow q$ $q \rightarrow r$ <hr/> Therefore, $p \rightarrow r$
Conjunction $p$ $q$ <hr/> Therefore, $p \wedge q$	Addition $p$ <hr/> Therefore, $p \vee q$	Resolution $p \vee q$ $\neg p \vee r$ <hr/> Therefore, $q \vee r$	Disjunctive Syllogism $p \vee q$ $\neg p$ <hr/> Therefore, $q$
Universal Instantiation $\forall x P(x)$ <hr/> Therefore, $P(c)$	Universal Generalization $P(c)$ <hr/> Therefore, $\forall x P(x)$	Existential Instantiation $\exists x P(x)$ <hr/> Therefore, $P(c)$	Existential Generalization $P(c)$ <hr/> Therefore, $\exists x P(x)$

## Inference Rules For Program Proofs

Composition Rule $p\{S_1\}q$ $q\{S_2\}r$ <hr/> $p\{S_1; S_2\}r$	Conditional Rule $(p \wedge \text{condition})\{S\}q$ $(p \wedge \neg \text{condition}) \rightarrow q$ <hr/> $p\{\text{if condition } S\}q$	Conditional with Else Rule $(p \wedge \text{condition})\{S_1\}q$ $(p \wedge \neg \text{condition})\{S_2\}q$ <hr/> $p\{\text{if condition } S_1 \text{ else } S_2\}q$
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