## Assertions, pre/post-

 conditions and invariantsAssertions: Section 4.2 in Savitch (p. 239)
Loop invariants: Section 4.5 in Rosen

## Programming as a contract

- Specifying what each method does
- Specify it in a comment before method's header
- Precondition
- What is assumed to be true before the method is executed
- Caller obligation
- Postcondition
- Specifies what will happen if the preconditions are met - what the method guarantees to the caller
- Method obligation

```
Example
** precondition: x >= 0
** postcondition: return value satisfies:
** result * result == x
*/
double sqrt(double x) {
}
```


## Class Invariants

- A class invariant is a condition that all objects of that class must satisfy while it can be observed by clients
- What about a Rectangle object?

What is an assertion?

- An assertion is a statement that says something about the state of your program
- Should be true if there are no mistakes in the program
$/ / \mathrm{n}=1$
while ( n < limit) \{
$\mathrm{n}=2$ * n ;
\}
// what could you state here?


## What is an assertion?

- An assertion is a statement that says something about the state of your program
- Should be true if there are no mistakes in the program

$$
/ / \mathrm{n}==1
$$

while ( n < limit) $\{$
$\mathrm{n}=2$ * n ;
\}
//n >= limit
assert
Using assert:
assert $\mathrm{n}==1$;
while (n < limit) \{ $\mathrm{n}=2$ * n ;
\}
assert n >= limit

```
When to use Assertions
- Another example
if (i % 3 == 0) { ... }
else if (i % 3 == 1) { ... }
else { // We know (i % 3 == 2)
... }
```

When to use Assertions

- We can use assertions to guarantee the behavior.
if (i \% 3 == 0) \{ ... \}
else if (i \% $3=1$ ) \{ ... \}
else \{ assert i \% $3=2$; ... \}

| Another example |
| :--- |
| int $\mathrm{p}=. ., \mathrm{d}=.$. ; |
| int $\mathrm{q}=\mathrm{p} / \mathrm{d} ;$ |
| int $\mathrm{r}=\mathrm{p} \% \mathrm{~d} ;$ |
| assert ? |

## Control Flow

- If a program should never reach a point, then a constant false assertion may be used

```
private void search() {
            for (...) {
                if (found) // will always happen
                return;
        }
        assert false; // should never get here
    }
```


## When to use assertions?

- Programming by contract
- Preconditions in methods (eg value ranges of parameters) should be enforced rather than asserted because assertions can be turned off
- Postconditions
- Assert post-condition


## Assertions in Eclipse

- To enable assert statements, you must set a compiler flag. Go to Run -> Run Configurations > Arguments, and in the box labeled VM arguments, enter either -enableassertions or just -ea


## Assertions

- Syntax:
assert Boolean_Expression;
- Each assertion is a Boolean expression that you claim is true.
- By verifying that the Boolean expression is indeed true, the assertion confirms your claims about the behavior of your program, increasing your confidence that the program is free of errors.
- If assertion is false when checked, the program raises an exception.


## Performance

- Assertions may slow down execution. For example, if an assertion checks to see if the element to be returned is the smallest element in the list, then the assertion would have to do the same amount of work that the method would have to do
- Therefore assertions can be enabled and disabled
- Assertions are, by default, disabled at run-time
- In this case, the assertion has the same semantics as an empty statement
- Think of assertions as a debugging tool
- Don't use assertions to flag user errors, because assertions can be turned off


## More Information

- For more information:
http://java.sun.com/j2se/1.4.2/docs/guide/ lang/assert.html


## Loop invariants

- We can use predicates (logical expressions) to reason about our programs.
- A loop invariant is a predicate
- that is true directly before the loop executes
- that is true before and after the loop body executes
- and that is true directly after the loop has executed
i.e., it is kept invariant by the loop.


## What does it mean...



## Loop invariants

- Combined with the loop condition, the loop invariant allows us to reason about the behavior of the loop:
<loop invariant>
while(test)\{ <test AND loop invariant> S;
<loop invariant>
\}
< not test AND loop invariant>

| <loop invariant> | If we can prove that |
| :---: | :---: |
| while(test)\{ | the loop invariant holds before the loop |
| <test AND | and that |
| loop invariant> S; | the loop body keeps the loop invariant true i.e. <test AND loop invariant> S; <loop invariant> |
| <loop invariant> | then we can infer that |
| , | not test AND loop invariant |
| < not test AND | holds after the loop terminates |

Example: loop index value after loop
<precondition: $\mathrm{n}>0$ >
int $\mathrm{i}=0$;
while $(\mathrm{i}<\mathrm{n})\{$
$i=i+1$;
\}
<post condition: $\mathrm{i}==\mathrm{n}>\quad$ We want to prove
$i==n$ right after the loop

Example: loop index value after loop

| // precondition: $\mathrm{n}>0$ |  |
| :---: | :---: |
| int $\mathrm{i}=0$; |  |
| // i<=n loop invariant while ( $\mathrm{i}<\mathrm{n}$ ) $\{$ |  |
| // $\mathrm{i}<\mathrm{n}$ test passed |  |
| // AND | So we can conclude the |
| // i<=n loop invariant | obvious: |
| i++; | $\mathrm{i}=$ = right after the loop |
|  |  |
| // i>=n AND $\mathrm{i}<=\mathrm{n} \rightarrow \mathrm{i}==\mathrm{n}$ |  |

## Example: sum of elements in an array

```
int total (int[] elements){
    int sum = 0, i = 0, n = elements.length;
    // sum == sum of elements from 0 to i-1
    while (i < n){
            // sum == sum of elements 0...i-1
            sum += elements [i];
            i++;
            // sum == sum of elements 0...i-1
    }
    // i==n (previous example) AND
    // sum == sum elements 0...i-1
    // }->\mathrm{ sum == sum of elements 0...n-1
    return sum;
}
```


## Closed Curve Game

- There are two players, Red and Blue. The game is played on a rectangular grid of points:

$$
\begin{aligned}
& \text { • . . . . . . } \\
& \text { •••••• } \\
& \text {. . . . . . . } \\
& \text {. . . . . . }
\end{aligned}
$$

Red draws a red line segment, either horizontal or vertical, connecting any two adjacent points on the grid that are not yet connected by a line segment. Blue takes a turn by doing the same thing, except that the line segment drawn is blue. Red's goal is to form a closed curve of red line segments. Blue's goal is to prevent Red from doing so

See

## Closed Curve Game

- Answer: Yes! Blue is guaranteed to win the game by responding to each turn by Red in the following manner:
if (Red drew a horizontal line segment) eet $i$ and $j$ such that Red's line segment connects ( $i, j$ ) with ( $i, j+1$ if (i>1)
$(i-1, j+1)$ with ( $i, j+1)$
\} else \{
draw a line segment anywhere
\} else // Red drew a vertical line segment
let $i$ and $j$ be such that Red's line segment connects ( $i, j$ ) with ( $i+1, j$ )
if ( $j>1$ ) \{
draw a horizontal line segment connecting (i+1,j-1) with (i+1,j)
else $\left\{\begin{array}{l}\text { draw a } \\ \text { draw } \\ \text { a line }\end{array}\right.$
draw a line segment anywhere
,
\}
$\qquad$


## Closed Curve Game

- We can express this game as a computer program:

```
while (more line segments can be drawn) {
            Red draws line segment;
            Blue draws line segment;
    }
```

Question: Does either Red or Blue have a winning strategy?

## Closed Curve Game

- By following this strategy Blue guarantees that Red does not have an "upper right corner" at any step.
- So, the invariant is:

There does not exist on the grid a pair of red line segments that form an upper right corner.

And in particular, Red has no closed curve!
$\qquad$

Can we show it works? Loop invariants!!

```
// pre: left >0 AND right >0
int a=left, b=right, p=0; //p: the product
// p + (a*b) == left * right loop invariant
while (a!=0){
    // a!=0 and p + (a*b) == left * right
    // loop condition and loop invariant
    if (odd(a)) p+=b;
    a/=2;
    b*=2;
    // p+(a*b) == left*right
}
// a==0 and p+a*b == left*right -> p == left*right
```



Try it on $8 * 7$
left right $a \quad b \quad p$
$\begin{array}{lllll}8 & 7 & 8 & 7 & 0\end{array}$
$414 \quad 0$
$2 \quad 28 \quad 0$
1560
0118 +=b: 56

| Relation to int representation $19 * 5$ |
| :---: |
| 00101 |
| 10011 |
| 1015 |
| 101010 |
| 00000 |
| 00000 |
| 101000080 |
| $\overline{1011111} 95$ |

Summary: Loop Invariant Reasoning
//loop invariant true before loop while (b) \{
// b AND loop invariant
S;
// loop invariant
\}
// not b AND loop invariant
not b helps you make a stronger observation than loop invariant alone.

