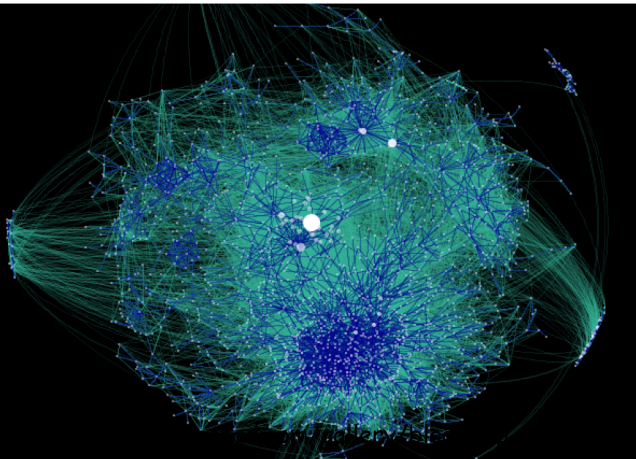


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# CS 220: Discrete Structures and their Applications

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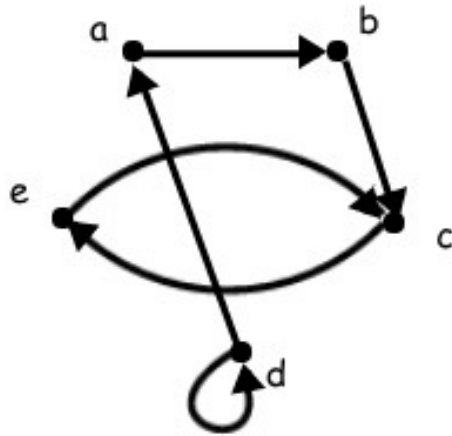
relations and directed  
graphs; transitive closure  
zybooks 9.3-9.6



# binary relations on a set

A binary relation on a set  $A$  is a subset of  $A \times A$ .

Graphical representation of a binary relation on a set:



self loop

$$A = \{a, b, c, d, e\}$$

$$R \subseteq A \times A$$

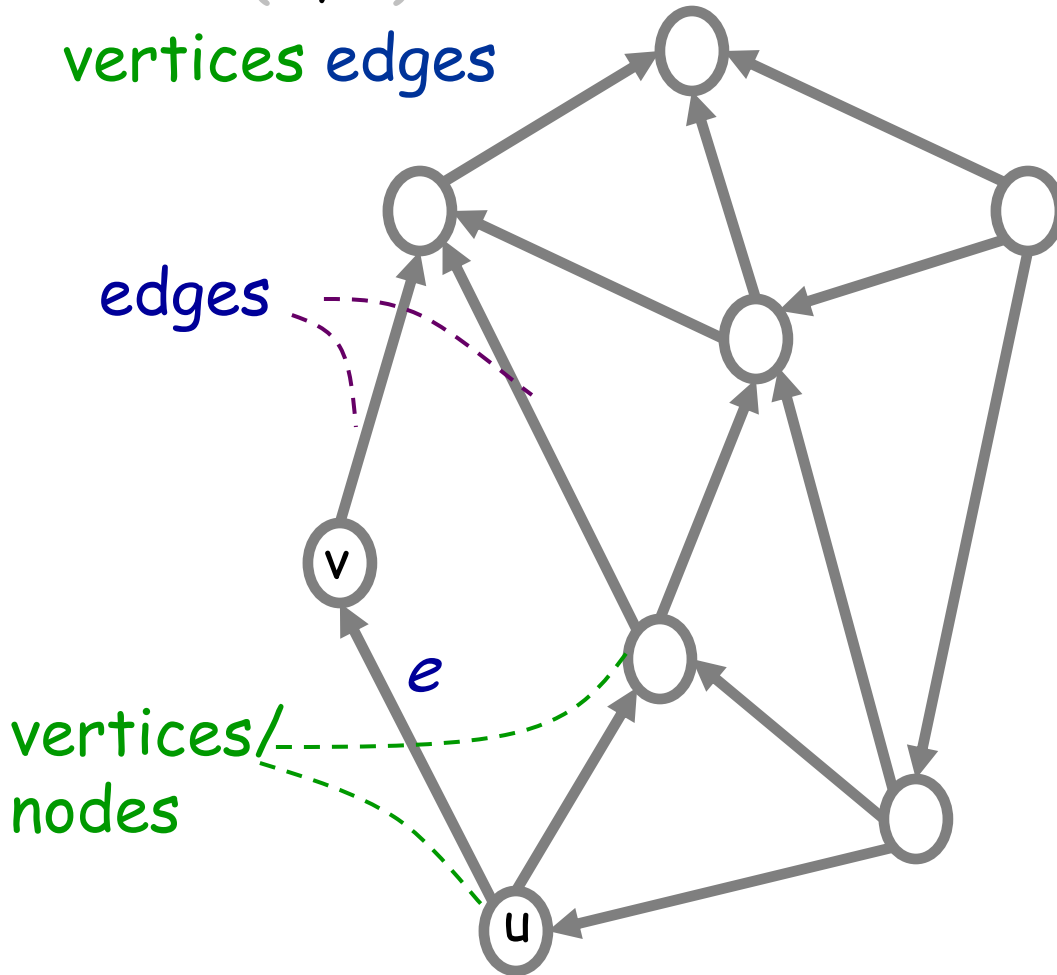
$$R = \{(a, b)(b, c)(e, c)(c, e)(d, a)(d, d)\}$$

This special case of a binary relation is also called a **directed graph**

# directed graphs

$$G=(V, E)$$

vertices edges



Edge  $(u, v)$  goes from vertex  $u$  to vertex  $v$ .

**in-degree** of a vertex: the number of edges pointing into it.

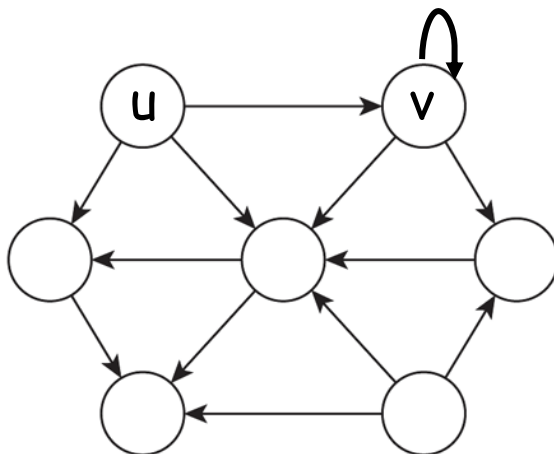
**out-degree** of a vertex: the number of edges pointing out of it.

# terminology

A directed graph (or digraph) is a pair  $(V, E)$ .

$V$  is a set of vertices, and  $E$ , a set of directed edges, is a subset of  $V \times V$ .

The vertex  $u$  is the **tail** of the edge  $(u, v)$  and vertex  $v$  is the **head**. If the head and the tail of an edge are the same vertex, the edge is called a **self-loop**.



Example: The web. What are the vertices/edges?

# matrices

An  $n \times m$  **matrix** over a set  $S$  is an array of elements from  $S$  with  $n$  rows and  $m$  columns.

The entry in row  $i$  and column  $j$  is denoted by  $A_{i,j}$ .

A matrix is called a **square** matrix if the number of rows is equal to the number of columns.

Is the adjacency matrix associated with a graph square?

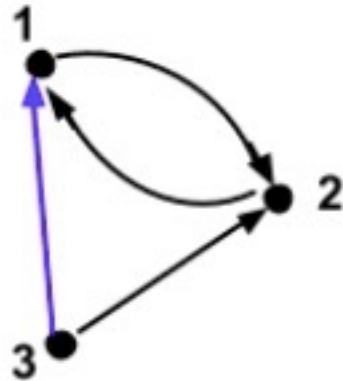
$$\begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$$

2 x 2 matrix  
over  $\{0, 1\}$

# adjacency matrix

A directed graph  $G$  with  $n$  vertices can be represented by an  $n \times n$  matrix over the set  $\{0, 1\}$  called the **adjacency matrix** for  $G$ .

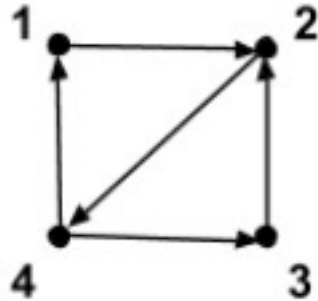
If  $A$  is the adjacency matrix for a graph  $G$ , then  $A_{i,j} = 1$  if there is an edge from vertex  $i$  to vertex  $j$  in  $G$ . Otherwise,  $A_{i,j} = 0$ .



$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{pmatrix}$$

# adjacency matrix

What are the missing values in the following adjacency matrix?



$$\begin{pmatrix} 0 & 1 & 0 & 0 \\ ? & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & ? & 0 \end{pmatrix}$$

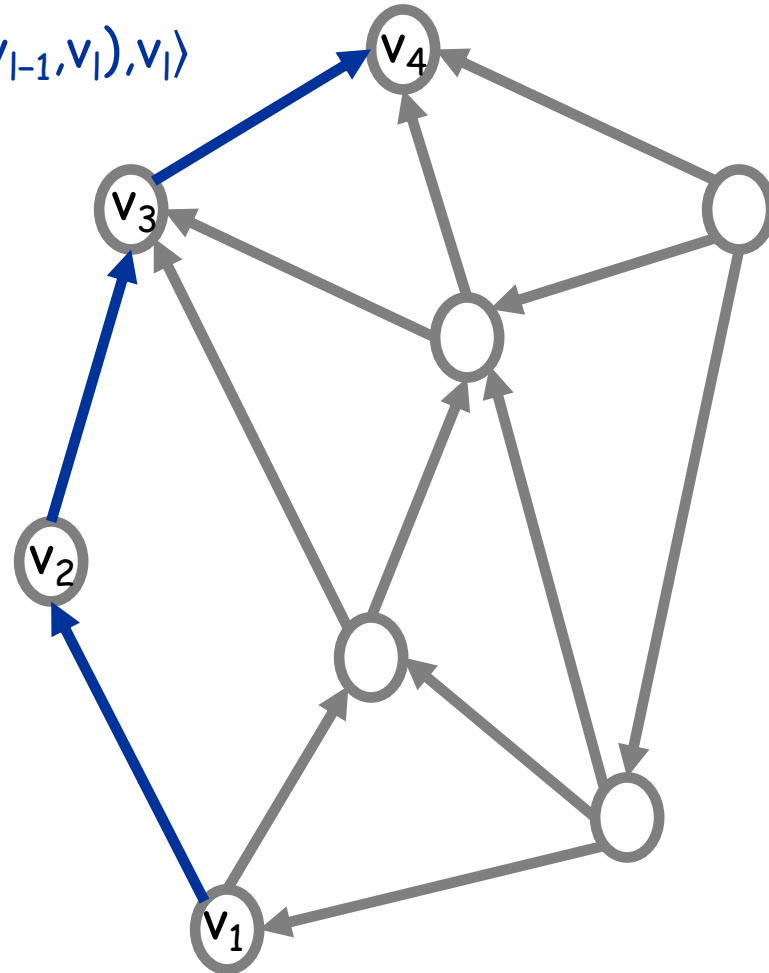
$$A_{2,1} = ?$$

$$A_{4,3} = ?$$

# walks

A **walk** from  $v_0$  to  $v_l$  in a directed graph  $G$  is a sequence of alternating vertices and edges that starts and ends with a vertex:

$\langle v_0, (v_0, v_1), v_1, (v_1, v_2), v_2, \dots, v_{l-1}, (v_{l-1}, v_l), v_l \rangle$





# walks

A **walk** from  $v_0$  to  $v_l$  in a directed graph  $G$  is a sequence of alternating vertices and edges that starts and ends with a vertex:

$$\langle v_0, (v_0, v_1), v_1, (v_1, v_2), v_2, \dots, v_{l-1}, (v_{l-1}, v_l), v_l \rangle$$

A walk can also be denoted by the sequence of vertices:

$$\langle v_0, v_1, \dots, v_l \rangle.$$

The sequence of vertices is a walk only if  $(v_{i-1}, v_i) \in E$  for  $i = 1, 2, \dots, l$ .

The **length** of a walk is  $l$ , the number of edges in the walk.

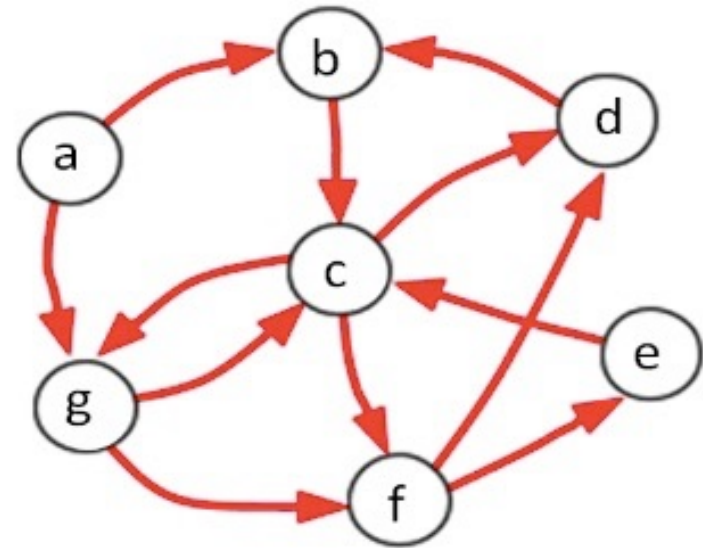
# walks, circuits, paths, cycles

A **circuit** is a walk in which the first vertex is the same as the last vertex.

A sequence of one vertex, denoted  $\langle a \rangle$ , is a circuit of length 0.

A walk is a **path** if no vertex is repeated in the walk.

A circuit is a **cycle** if there are no other repeated vertices, except the first and the last.



# composite relations

Let  $R$  be a relation from  $A$  to  $B$ , and let  $S$  be a relation from  $B$  to  $C$ . The **composite**  $S \circ R$  of  $R$  and  $S$  is defined as:

$$S \circ R = \{(a,c) : \exists b \text{ such that } aRb \text{ and } bSc\}$$

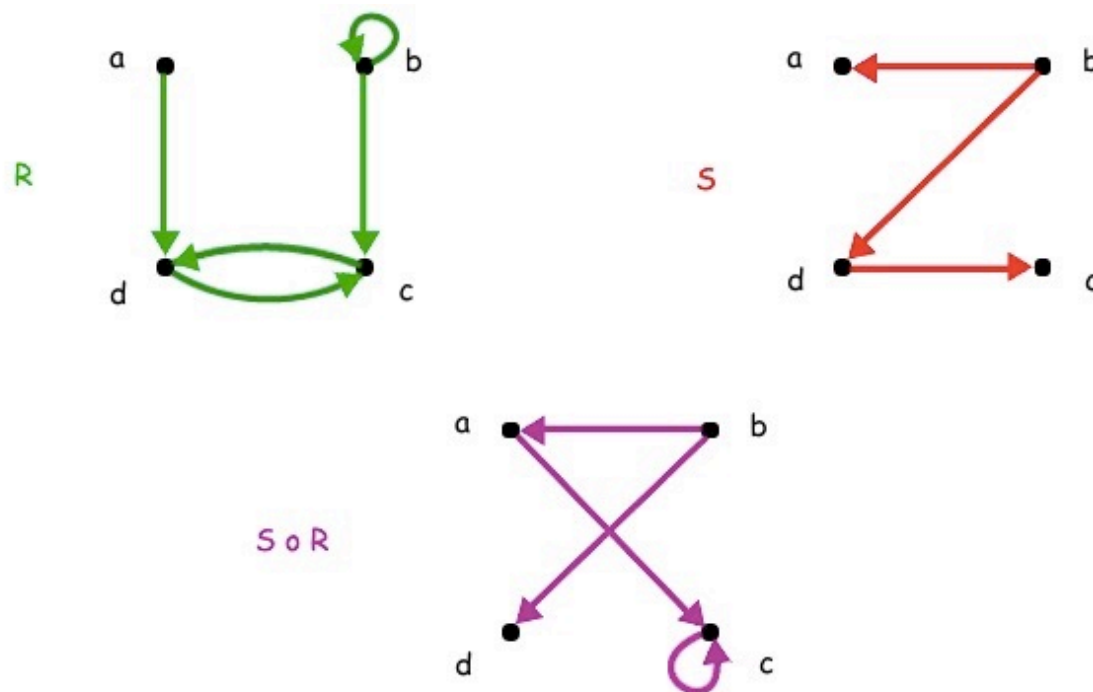
**Example:** Let  $R$  be the relation such that  $aRb$  if  $a$  is a parent of  $b$ . What is the relation  $R \circ R$ ?

# composite relations

Let  $R$  be a relation from  $A$  to  $B$ , and let  $S$  be a relation from  $B$  to  $C$ . The **composite**  $S \circ R$  of  $R$  and  $S$  is defined as:

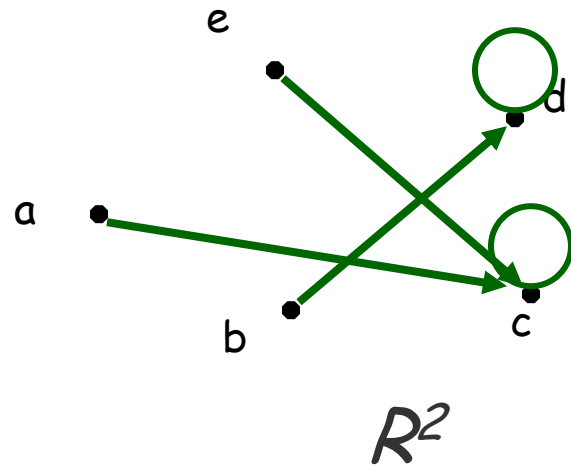
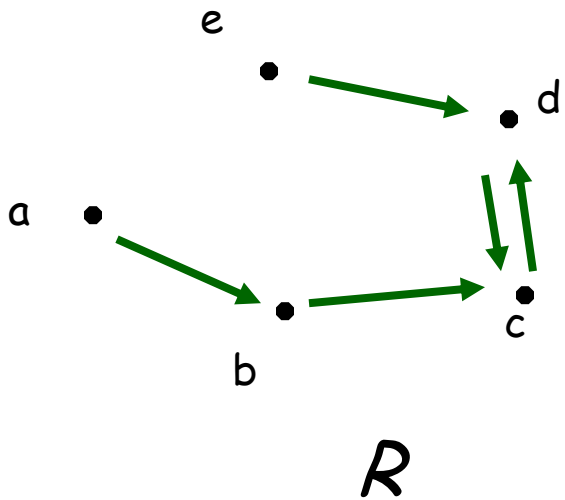
$$S \circ R = \{(a,c) : \exists b \text{ such that } aRb \text{ and } bSc\}$$

Example:



# composite relations

Composite relation on a set:



# composite relations

The powers  $R^n$  of relation  $R$  can be defined recursively:

$$R^1 = R \text{ and } R^{n+1} = R^n \circ R$$

The statement of *six-degrees of separation* can be succinctly expressed as  $aR^6b$  for all  $a, b$  where  $R$  is the relation on the set of people such that  $aRb$  if  $a$  knows  $b$

# paths and relations

The edge set  $E$  of a directed graph  $G$  can be viewed as a relation.

$E^k$  : the relation  $E$  composed with itself  $k$  times.

$G^k$  : the directed graph whose edge set is  $E^k$ .

**The Graph Power Theorem:** Let  $G$  be a directed graph. Let  $u$  and  $v$  be any two vertices in  $G$ . There is an edge from  $u$  to  $v$  in  $G^k$  if and only if there is a walk of length  $k$  from  $u$  to  $v$  in  $G$ .

# paths and relations

**The Graph Power Theorem:** Let  $G$  be a directed graph. Let  $u$  and  $v$  be any two vertices in  $G$ . There is an edge from  $u$  to  $v$  in  $G^k$  if and only if there is a walk of length  $k$  from  $u$  to  $v$  in  $G$ .

Proof by induction on  $k$ .

Base case: there is a walk of length 1 iff  $aRb$

Induction step: assume that there is an edge from  $u$  to  $v$  in  $G^k$  if and only if there is a walk of length  $k$  from  $u$  to  $v$  in  $G$ .

Let's prove this for  $k+1$ :

There is a walk of length  $k+1$  from  $a$  to  $b$  iff there is a  $c$  in  $A$  such that there is a walk of length 1 from  $a$  to  $c$ , i.e.  $aRc$  and a path of length  $k$  from  $c$  to  $b$ ; by the induction hypothesis this happens iff  $cR^k b$ , and by definition of

composition iff  $aR^{k+1}b$



# the transitive closure

The **transitive closure** of a graph  $G$ :

$$G^+ = G^1 \cup G^2 \cup G^3 \cup G^4 \dots$$

In the union, there is only one copy of the vertex set and the union is taken over the edge sets of the graphs.

$(u, v)$  is an edge in  $G^+$  if vertex  $v$  can be reached from vertex  $u$  in  $G$  by a walk of any length.

Similarly we can define the transitive closure of a relation  $R$ :

$$R^+ = R^1 \cup R^2 \cup R^3 \cup R^4 \dots$$

# the transitive closure

## Examples:

Let  $R$  be the relation between states in the US where  $aRb$  if  $a$  and  $b$  share a common border. What is  $R^+$ ?

What is  $R^+$  for the parent relation?

$A = \{1, 2, b\}$ . What is the transitive closure for:

- $R = \{(1, 1), (b, b)\}$
- $S = \{(1, 2), (2, b)\}$
- $T = \{(2, 1), (b, 2), (1, 1)\}$

# the transitive closure

The **transitive closure** of a graph  $G$ :

$$G^+ = G^1 \cup G^2 \cup G^3 \cup G^4 \dots$$

If the graph has  $n$  vertices:

$$G^+ = G^1 \cup G^2 \cup G^3 \cup \dots \cup G^n$$

The same holds for a relation  $R$ . Let  $R$  be a relation on a finite domain with  $n$  elements. Then

$$R^+ = R^1 \cup R^2 \cup R^3 \cup \dots \cup R^n$$

# the transitive closure

**Lemma:** Let  $G$  be a graph with  $n$  vertices. If there is a path from  $u$  to  $v$  in  $G$ , then there is such a path with length not exceeding  $n$ .

# an algorithm for the transitive closure

Let  $R$  be a relation over a set  $A$ .

Repeat the following step until no pair is added to  $R$ :

- ✓ If there are  $x, y, z \in A$  such that  $(x, y) \in R$ ,  $(y, z) \in R$  and  $(x, z) \notin R$ , then add  $(x, z)$  to  $R$ .