Frequently asked questions from the previous class survey

- Return type for `collect()`? Can we collect somewhere else?
- My team has 3 members: that means each of us can contribute a third less than a 2 member team, correct?
  - No such luck. In fact, the project is expected to be commensurately more complex.
Topics covered in this lecture

- Pair RDDs
- Dependencies and Transformations
- Partitioners
Transformations on Pair RDDs

1. Pair RDD = \{(1,2), (3,4), (3,6) \}

2. **reduceByKey(func)**
   - Combine values with the same key
   - Invocation: `rdd.reduceByKey((x, y) => x + y)`
   - Result: \{(1, 2), (3, 10) \}

3. **groupByKey(func)**
   - Group values with the same key
   - Invocation: `rdd.groupByKey()`
   - Result: \{(1, [2]), (3, [4, 6]) \}
Transformations on Pair RDDs [3/5]

- Pair RDD = \{(1,2), (3,4), (3,6) \}

- `mapValues(func)`
  - Apply function to each value of a pair RDD *without* changing the key
  - Invocation: `rdd.mapValues(x => x+1)`
  - Result: \{(1, 3), (3, 5), (3, 7) \}

Transformations on Pair RDDs [4/5]

- Pair RDD = \{(1,2), (3,4), (3,6) \}

- `values()`
  - Return an RDD of just the values
  - Invocation: `rdd.values()`
  - Result: \{2, 4, 6 \}
Transformations on Pair RDDs

- Pair RDD = {(1,2), (3,4), (3,6) }

- `sortByKey()`
  - Return an RDD sorted by the key
  - Invocation: `rdd.sortByKey()`
  - Result: { (1,2), (3,4), (3,6) }

**Transformations on two pair RDDs**
**Transformations on two Pair RDDs**

- \[ \text{rdd} = \{(1,2), (3,4), (3,6)\} \quad \text{other} = \{(3,9)\} \]

- `subtractByKey()`
  - Remove elements with a key present in the `other` RDD
  - Invocation: `rdd.subtractByKey(other)`
  - Result: \{(1,2)\}

- `join()`
  - Perform an **inner join** between two RDDs. Only keys that are present in both pair RDDs are output
  - Invocation: `rdd.join(other)`
  - Result: \{(3, (4,9)), (3, (6,9))\}
Transformations on two Pair RDDs [3/5]

- **rdd** = {{1,2}, (3,4), (3,6)}    **other** = {{3,9}}

- **leftOuterJoin()**
  - Perform a join between two RDDs where the key must be present in the first RDD
  - Value associated with each key is a tuple of the value from the source and an Option for the value from the other pair RDD
    - In python if a value is not present, None is used.
  - Invocation: `rdd.leftOuterJoin(other)`
  - Result: { (1, (2,None)) , (3, (4, 9)) , (3, (6, 9)) }

Transformations on two Pair RDDs [4/5]

- **rdd** = {{1,2}, (3,4), (3,6)}    **other** = {{3,9}}

- **rightOuterJoin()**
  - Perform a join between two RDDs where the key must be present in the other RDD;
  - Tuple has an option for the source rather than other RDD
  - Invocation: `rdd.rightOuterJoin(other)`
  - Result: { (3, (4,9) ), (3, (6,9)) }
Transformations on two Pair RDDs

- \( \text{rdd} = \{(1,2), (3,4), (3,6)\} \quad \text{other} = \{(3,9)\} \)

- **cogroup()**
  - Group data from both RDDs using the same key
  - Invocation: \( \text{rdd}.\text{cogroup}(\text{other}) \)
  - Result: \( \{(1, ([2], [])), (3, ([4, 6], [9]))\} \)

Example of chaining operations:
Calculation of per-key average

<table>
<thead>
<tr>
<th>key</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>panda</td>
<td>0</td>
</tr>
<tr>
<td></td>
<td>3</td>
</tr>
<tr>
<td></td>
<td>3</td>
</tr>
<tr>
<td></td>
<td>1</td>
</tr>
<tr>
<td></td>
<td>4</td>
</tr>
</tbody>
</table>

```
rdd.mapValues(x=> (x, 1)).reduceByKey( (x, y) => (x._1 + y._1, x._2 + y._2))
```
A word count example

- We are using `flatMap()` to produce a pair RDD of words and the number 1

```python
rdd = sc.textfile("s3://...")
words = rdd.flatMap(lambda x: x.split(" "))
result = words.map(lambda x: (x,1)).reduceByKey(lambda x, y: x+y)
```

Wide and Narrow Transformations

October 3, 2017
Transformations and Dependencies

- Two categories of dependencies
  - Narrow
    - Each partition of the parent RDD is used by at most one partition of the child RDD
  - Wide
    - Multiple child RDD partitions may depend on a single parent RDD partition

- The narrow versus wide distinction has significant implications for the way Spark evaluates a transformation and, consequently, for its performance

Narrow Transformations

- Narrow transformations are those in which each partition in the child RDD has simple, finite dependencies on partitions in the parent RDD
- Can be determined at design time, irrespective of the values of the records in the parent partitions
- Partitions in narrow transformations can either depend on:
  - One parent (such as in the map operator), or
  - A unique subset of the parent partitions that is known at design time (coalesce)
- Narrow transformations can be executed on an arbitrary subset of the data without any information about the other partitions.
Dependencies between partitions for narrow transformations

Wide Transformations

- Transformations with **wide dependencies** cannot be executed on arbitrary rows
- Require the data to be partitioned in a particular way, e.g., according the value of their key
  - In sort, for example, records have to be partitioned so that keys in the same range are on the same partition
- Transformations with wide dependencies include sort, reduceByKey, groupByKey, join, and anything that calls the rePartition function
Dependencies between partitions for wide transformations

Wide dependencies cannot be known fully before the data is evaluated.

The dependency graph for any operations that cause a shuffle (such as groupByKey, reduceByKey, sort, and sortByKey) follows this pattern.

Tuning the level of parallelism
Tuning the level of parallelism

- Every RDD has a **fixed number of partitions**
  - Determine the degree of parallelism when executing operations

- During aggregations or grouping operations, you can ask Spark to use a specific number of partitions
  - This will override defaults that Spark uses

---

Example: Tuning the level of parallelism

```python
data = ["a", 3], ("b", 4), ("a", 1)]

sc.parallelize(data).
    reduceByKey(lambda x, y: x + y) #default

sc.parallelize(data).
    reduceByKey(lambda x, y: x + y, 10) #Custom
```
What if you want to tune parallelism outside of grouping and aggregation operations?

- There is repartition()
  - Shuffles data across the network to create a new set of partitions
  - Very expensive operation!

- There is the coalesce() operation
  - Allows avoiding data movement
    - But only if you are decreasing the number of partitions
  - Check rdd.getNumPartitions() and make sure you are coalescing to fewer partitions than current

PAIR RDDs: What to watch for
Despite their utility, key/value operations can lead to a number of performance issues

- Most expensive operations in Spark fit into the key/value pair paradigm
  - Because *most wide transformations* are key/value transformations,
  - And most require some fine tuning and care to be performant

In particular, operations on key/value pairs can cause

1. Out-of-memory errors in the driver
2. Out-of-memory errors on the executor nodes
3. Shuffle failures
4. “Straggler tasks” or partitions, which are especially slow to compute

- The last three performance issues are all most often caused by *shuffles associated with the wide transformations*
Memory errors in the driver, is usually caused by actions

- Several key/value actions (including countByKey, countByValue, lookUp, and collectAsMap) return data to the driver.
- In most instances they return unbounded data since the number of keys and the number of values are unknown.
- In addition to number of records, the size of each record is an important factor in causing memory errors.

Preventing out-of-memory errors with aggregation operations [1/2]

- combineByKey and all of the aggregation operators built on top of it (reduceByKey, foldLeft, foldRight, aggregateByKey) may lead to memory errors if they cause the accumulator to become too large for one key.
- What about groupByKey?
  - It is actually implemented using combineByKey where the accumulator is an iterator with all the data.
Preventing out-of-memory errors with aggregation operations [2/2]

- Use functions that implement **map-side combinations**
  - Meaning that records with the same key are combined before they are shuffled
  - This can greatly reduce the shuffled read

- The following four functions are implemented to use map-side combinations
  - reduceByKey
  - treeAggregate
  - aggregateByKey
  - foldByKey

Two primary techniques to avoid performance problems associated with shuffles

- Shuffle Less
- Shuffle Better
Shuffle Less

- Preserve partitioning across narrow transformations to avoid reshuffling data
- Use the same partitioner on a sequence of wide transformations. This can be particularly useful:
  - To avoid shuffles during joins and ...
  - To reduce the number of shuffles required to compute a sequence of wide transformations

Shuffle Better

- Sometimes, computation cannot be completed without a shuffle
- However, not all wide transformations and not all shuffles are equally expensive or prone to failure
Shuffle Better [2/2]

- By using wide transformations such as `reduceByKey` and `aggregateByKey` that can perform map-side reductions and that do not require loading all the records for one key into memory?
  - You can prevent memory errors on the executors and
  - Speed up wide transformations, particularly for aggregation operations

- Lastly, shuffling data in which records are distributed evenly throughout the keys, and which contain a high number of distinct keys?
  - Prevents out-of-memory errors on the executors and “straggler tasks”
Partitioners

- The partitioner defines **how records will be distributed** and thus which records will be completed by each task.

- Practically, a partitioner is actually an interface with two methods:
  - `numPartitions` that defines the number of partitions in the RDD after partitioning.
  - `getPartition` that defines a mapping from a key to the integer index of the partition where records with that key should be sent.

There are two implementations for the partitioner object provided by Spark:

- **HashPartitioner**
  - Determines the index of the child partition based on the hash value of the key.

- **RangePartitioner**
  - Assigns records whose keys are in the same range to a given partition.
  - Required for sorting since it ensures that by sorting records within a given partition, the entire RDD will be sorted.

- It is possible to define a custom partitioner.
Partitioners and transformations

- Unless a transformation is known to only change the value part of the key/value pair in Spark
  - The resulting RDD will **not have a known** partitioner
  - Even if the partitioning has not changed

Using narrow transformations that preserve partitioning

- Some narrow transformations, such as `mapValues`, **preserve the partitioning** of an RDD if it exists
- Common transformations like `map` and `flatMap` can change the key
  - So even if your function does not change the key, the resulting RDD will not have a known partitioner.
  - Instead, if you don’t want to modify the keys, call the `mapValues` function (defined only on pair RDDs)
    - It keeps the keys, and therefore the partitioner, exactly the same.
    - The `mapPartitions` function will also preserve the partition if the `preservesPartitioning` flag is set to true.
Spark Streaming

- Act on data **as soon as it arrives**
  - Track statistics of page views in real time, detect anomalies, etc.

- Spark streaming
  - Spark’s module for dealing with streaming data
  - Uses an API very similar to what we have seen with batch jobs (centered around RDDs)

- Available in Java and Scala
  - Recent support for Python
Spark Streaming: Core concepts

- Provides an abstraction called DStreams (discretized streams)
- A DStream is a sequence of data arriving over time
- Internally, a DStream is represented as a sequence of RDDs arriving at each time step

DStreams

- DStreams can be created from various input sources
  - Flume, Kafka, or HDFS
- Once built, DStreams offer two types of operations:
  - Transformations: Yields a new DStream
  - Output operations: Writes data to an external system
- Provides many of the same operations available on RDDs
  - PLUS new operations related to time (e.g. sliding windows)
Example

- Start by creating a `StreamingContext`
  - Main entry point for streaming functionality
  - Specify batch interval, specifying **how often** to process new data

- We will use `socketTextStream()` to create a DStream based on text data received over a port

- Transform DStream with filter to get lines that contain “error”

```
JavaStreamingContext jssc =
    new JavaStreamingContext(conf, Durations.seconds(1));

JavaDStream<String> lines =
    jssc.socketTextStream("localhost", 7777);

JavaDStream<String> errorLines =
    lines.filter(new Function<String, Boolean>() {
        public Boolean call(String line) {
            return line.contains("error");
        }
    });
```
Previous snippet only sets up the computation

- To start receiving the data?
  - Explicitly call start() on StreamContext

- SparkStreaming will start to schedule Spark jobs on the underlying SparkContext
  - Occurs in a separate thread
  - To keep application from terminating?
    - Also call awaitTermination()

```java
jssc.start();
jssc.awaitTermination();
```

The contents of this slide-set are based on the following references