

## Compiling for Parallelism & Locality

### Assignments

- Deadline for project 4 extended to Dec 1

### Last time

- Data dependences and loops

### Today

- Finish data dependence analysis for loops

## Dependence Testing in General

### General code

```
do i1 = l1, h1
  ...
  do in = ln, hn
    A(f(i1, ..., in)) = ... A(g(i1, ..., in))
  enddo
  ...
enddo
```

There exists a dependence between iterations  $I=(i_1, \dots, i_n)$  and  $J=(j_1, \dots, j_n)$  when

- $f(I) = g(J)$
- $(l_1, \dots, l_n) < I, J < (h_1, \dots, h_n)$

## Algorithms for Solving the Dependence Problem

### Heuristics

- GCD test (Banerjee76, Towle76): determines whether integer solution is possible, no bounds checking
- Banerjee test (Banerjee 79): checks real bounds
- I-Test (Kong et al. 90): integer solution in real bounds
- Lambda test (Li et al. 90): all dimensions simultaneously
- Delta test (Goff et al. 91): pattern matches for efficiency
- Power test (Wolfe et al. 92): extended GCD and Fourier Motzkin combination

Use some form of Fourier-Motzkin elimination for integers, exponential worst-case

- Parametric Integer Programming (Feautrier91)
- Omega test (Pugh92)

## Dependence Testing

Consider the following code...

```
do i = 1, 5
  A(3*i+2) = A(2*i+1)+1
enddo
```

### Question

- How do we determine whether one array reference depends on another across iterations of an iteration space?

## Dependence Testing: Simple Case

### Sample code

```
do i = 1,h
  A(a*i+c1) = ... A(a*i+c2)
enddo
```

### Dependence?

- $a*i_1+c_1 = a*i_2+c_2$ , or
- $a*i_1 - a*i_2 = c_2-c_1$
- Solution exists if  $a$  divides  $c_2-c_1$

## Example

### Code

```
do i = 1,h
  A(2*i+2) = A(2*i-2)+1
enddo
```

$i_1$

$i_2$

### Dependence?

$2*i_1 - 2*i_2 = -2 - 2 = -4$   
(yes, 2 divides -4)

### Kind of dependence?

- Anti?  $i_2 + d = i_1 \Rightarrow d = -2$
- Flow?  $i_1 + d = i_2 \Rightarrow d = 2$

## GCD Test

### Idea

- Generalize test to linear functions of iterators

### Code

```
do i = 1i,hi
  do j = 1j,hj
    A(a1*i + a2*j + a0) = ... A(b1*i + b2*j + b0) ...
  enddo
enddo
```

### Again

- $a_1*i_1 - b_1*i_2 + a_2*j_1 - b_2*j_2 = b_0 - a_0$
- Solution exists if  $\text{gcd}(a_1, a_2, b_1, b_2)$  divides  $b_0 - a_0$

## Example

### Code

```
do i = 1i,hi
  do j = 1j,hj
    A(4*i + 2*j + 1) = ... A(6*i + 2*j + 4) ...
  enddo
enddo
```

$\text{gcd}(4, -6, 2, -2) = 2$

Does 2 divide 4-1?

## Banerjee Test

```
for (i=L; i<=U; i++) {
  x[a_0 + a_1*i] = ...
  ... = x[b_0 + b_1*i]
}
```

Does  $a_0 + a_1*i = b_0 + b_1*i'$  for some integer  $i$  and  $i'$ ?  
 If so then  $(a_1*i - b_1*i') = (b_0 - a_0)$

Determine upper and lower bounds on  $(a_1*i - b_1*i')$

```
for (i=1; i<=5; i++) {
  x[i+5] = x[i];
}
```

upper bound =  $a_1*\max(i) - b_1*\min(i') = 4$   
 lower bound =  $a_1*\min(i) - b_1*\max(i') = -4$   
 $b_0 - a_0 =$

## Distance Vectors: Legality

### Definition

– A dependence vector,  $v$ , is **lexicographically nonnegative** when the left-most entry in  $v$  is positive or all elements of  $v$  are zero

Yes:  $(0,0,0)$ ,  $(0,1)$ ,  $(0,2,-2)$

No:  $(-1)$ ,  $(0,-2)$ ,  $(0,-1,1)$

– A dependence vector is **legal** when it is lexicographically nonnegative (assuming that indices increase as we iterate)

### Why are lexicographically negative distance vectors illegal?

### What are legal direction vectors?

## Direction Vector

### Definition

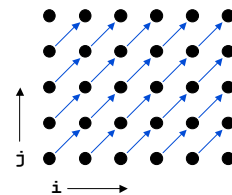
- A direction vector serves the same purpose as a distance vector when less precision is required or available
- Element  $i$  of a direction vector is  $<$ ,  $>$ , or  $=$  based on whether the source of the dependence precedes, follows or is in the same iteration as the target in loop  $i$

### Example

```
do i = 1, 6
  do j = 1, 5
    A(i, j) = A(i-1, j-1) + 1
  enddo
enddo
```

Direction vector:  $(<, <)$

Distance vector:  $(1, 1)$



## Loop-Carried Dependences

### Definition

– A dependence  $D=(d_1, \dots, d_n)$  is **carried** at loop level  $i$  if  $d_i$  is the first nonzero element of  $D$

### Example

```
do i = 1, 6
  do j = 1, 6
    A(i, j) = B(i-1, j) + 1
    B(i, j) = A(i, j-1) * 2
  enddo
enddo
```

Distance vectors:  $(0, 1)$  for accesses to **A**  
 $(1, 0)$  for accesses to **B**

### Loop-carried dependences

- The  $j$  loop carries dependence due to **A**
- The  $i$  loop carries dependence due to **B**

## Parallelization

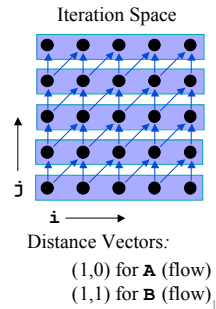
### Idea

- Each iteration of a loop may be executed in parallel if it carries no dependences

### Example

```
do i = 1, 6
  do j = 1, 5
    A(i, j) = B(i-1, j-1) + 1
    B(i, j) = A(i, j-1) * 2
  enddo
enddo
```

### Parallelize i loop?



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Data Dependence Analysis

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## Parallelization

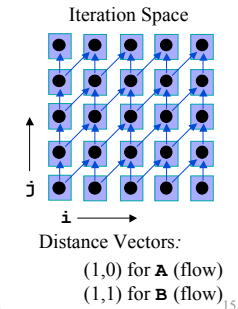
### Idea

- Each iteration of a loop may be executed in parallel if it carries no dependences

### Example

```
do i = 1, 6
  do j = 1, 5
    A(i, j) = B(i-1, j-1) + 1
    B(i, j) = A(i, j-1) * 2
  enddo
enddo
```

### Parallelize j loop?



CS553 Lecture

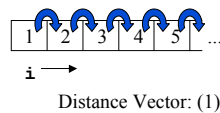
Data Dependence Analysis

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## Example 2: Parallelization (reprise)

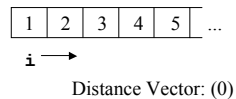
### Why can't this loop be parallelized?

```
do i = 1, 100
  A(i) = A(i-1) + 1
enddo
```



### Why can this loop be parallelized?

```
do i = 1, 100
  A(i) = A(i) + 1
enddo
```



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## Example 1: Loop Permutation (reprise)

### Sample code

```
do j = 1, 6
  do i = 1, 5
    A(j, i) = A(j, i) + 1
  enddo
enddo
```

↔

```
do i = 1, 5
  do j = 1, 6
    A(j, i) = A(j, i) + 1
  enddo
enddo
```

### Why is this legal?

- No loop-carried dependences, so we can arbitrarily change order of iteration execution

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## Concepts

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### Improve performance by ...

- improving data locality
- parallizing the computation

### Data Dependences

- iteration space
- distance vectors and direction vectors
- loop carried

### Transformation legality

- must respect data dependences
- scalar expansion as a technique to remove anti and output dependences

### Data Dependence Testing

- general formulation of the problem
- GCD test

## Next Time

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### Lecture

- Loop transformations for parallelism and locality