

High-Performance Embedded Systems-on-a-Chip

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Lecture 8: Review & Exercises

Prove closure of SURE's under same CoB

Consider the following SRE:

$$X[z] = \{z \in D_X\} : g(\dots, X[f_{xx}(z)], Y[f_{xy}(z)])$$

$$Y[z] = \{z \in D_Y\} : h(\dots, X[f_{yx}(z)], Y[f_{yy}(z)])$$

and a bijective affine function, $\mathcal{T} : z \mapsto Tz + t$ where T is an integer unimodular matrix, and t is a vector. Show that when all the dependence functions are **uniform** (i.e., the SRE is actually a SURE) applying the **same** transformation, \mathcal{T} , to **both** X and Y , the resulting system remains uniform. Determine the new dependence vectors. Can you think of a (slightly) more general transformation that retains this closure property?

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Applying CoB by \mathcal{T} first to X

$$X[z] = \{z \in \mathcal{T}(D_X)\} : g(\dots, X[\mathcal{T} \circ f_{xx} \circ \mathcal{T}^{-1}(z)], \\ Y[f_{xy} \circ \mathcal{T}^{-1}(z)])$$

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$$Y[z] = \{z \in \mathcal{T}(D_Y)\} : h(\dots, X[\mathcal{T} \circ f_{yx} \circ \mathcal{T}^{-1}(z)], \\ Y[\mathcal{T} \circ f_{yy} \circ \mathcal{T}^{-1}(z)])$$

Each new dependence is of the form $\mathcal{T} \circ f \circ \mathcal{T}^{-1}(z)$, where f is uniform, i.e., $f(z) = z + \delta$

$$\begin{aligned}\mathcal{T} \circ f \circ \mathcal{T}^{-1}(z) &= T f(T^{-1}z - T^{-1}t) + t \\ &= T(T^{-1}z - T^{-1}t + \delta) + t \\ &= (z - t + T\delta) + t \\ &= z + T\delta\end{aligned}$$

Bubble Sort

Starting SURE(subscripts on lhs for compaction)

$$\begin{aligned}
 y_i &= \begin{cases} i = 1 : m[n, n - 1] \\ i > 1 : M[n, n - i + 1] \end{cases} \\
 m_{i,j} &= \begin{cases} i = 2; j = 1 & : \min(x_1, x_2) \\ 2 < i \leq n; j = 1 & : \min(M[i - 1, j], x_i) \\ 3 \leq j + 1 = i \leq n & : \min[m[i - 1, j - 1], m[i, j - 1]] \\ j + 2 \leq i \leq n; 1 < j & : \min(M[i - 1, j], m[i, j - 1]) \end{cases} \\
 M_{i,j} &= \begin{cases} i = 2; j = 1 & : \max(x_1, x_2) \\ 2 < i \leq n; j = 1 & : \max(M[i - 1, j], x_i) \\ 3 \leq j + 1 = i \leq n & : \max[m[i - 1, j - 1], m[i, j - 1]] \\ j + 2 \leq i \leq n; 1 < j & : \max(M[i - 1, j], m[i, j - 1]) \end{cases}
 \end{aligned}$$

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$$\begin{pmatrix} t \\ p \end{pmatrix} = \mathcal{T} \begin{pmatrix} i \\ j \end{pmatrix} = \begin{pmatrix} i + j \\ j \end{pmatrix} = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} \begin{pmatrix} i \\ j \end{pmatrix}$$

Transformed SURE

$$y_i = \begin{cases} i = 1 : m[2n - 1, n - 1] \\ i > 1 : M[2n - i + 1, n - i + 1] \end{cases}$$

$$m_{t,p} = \begin{cases} t = 3; p = 1 & : \min(x_1, x_2) \\ 3 < t \leq n + 1; p = 1 & : \min(M[t - 1, p], x_{t-1}) \\ 5 \leq t = 2p + 1 \leq n + p & : \\ & \min[m[t - 2, p - 1], m[t - 1, p - 1]] \\ 2p + 2 \leq t \leq n + p; 1 < p & : \\ & \min(M[t - 1, p], m[t - 1, p - 1]) \end{cases}$$

$$M_{t,p} = \text{similar}$$

I/O and Control

Processors involved in I/O operations:

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Basic Rule: I/O must occur at the **boundary processors**