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Technical Report*



**ANTIRANDOM TESTING:
GETTING THE MOST OUT OF BLACK-BOX TESTING***

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ABSTRACT

Random testing is a well known concept that requires that each test is selected randomly regardless of the test previously applied. This paper introduces the concept of antirandom testing. In this testing strategy each test applied is chosen such that its total distance from all previous tests is maximum. Two distance measures are defined. Procedures to construct antirandom sequences are developed. A checkpoint encoding scheme is introduced that allows automatic generation of efficient test cases. Further developments and studies needed are identified.

1 Introduction

Exhaustive testing of software is infeasible except for very small programs [1, 2]. Achieving 100% test coverage using a specific measure does not assure that all defects have been found [4, 5]. Obtaining total coverage itself may be hard; 85% branch coverage is often used as a target. The testing time, often a significant fraction of the overall development time, is always limited. The testers thus have the challenging task of making testing as efficient as possible, since the cost of remaining defects in the released code can be very high.

Here we consider black-box testing where only the external specifications are used to obtain test suites. No implementation specific information is assumed to be known. Often software testing is termed *random* [3, 6, 7]. By definition, in random testing the values of the inputs in each test are selected randomly, regardless of the previous tests applied. Random testing and its variations have been extensively used and studied for hardware systems.

Random testing avoids the problem of deterministic test generation that requires structural information about the program to be processed for generating each test. Available evidence suggests that random testing may be a reasonable choice for obtaining a moderate degree of confidence; however, it becomes very inefficient when the residual defect density becomes low [6].

Random testing does not exploit some information that is available in black-box testing environment. This information consists of the previous tests applied. If an experienced tester is generating tests by hand, he would select each new test such that it covers some part of the functionality not yet covered by tests already generated. The objective of this paper is to formally define an approach that uses this information and to propose schemes that may allow such test generation to be done automatically. We term this approach *antirandom* testing, since selection of each test explicitly depends on the tests already obtained. The problem of generating antirandom sequences is first considered for boolean inputs. to make each new test as different as possible, we use *Hamming distance* and *Cartesian distance* as measures of difference. In general, the input variables for a program can be numbers, characters as well as data structures. Here we also present an approach to efficiently encode the input space spanned by such variables into binary. This allows binary antirandom sequences to be decoded into actual inputs.

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Test data selection has been regarded as an important part of testing software [8, 9, 10, 11]. Researchers have identified valuable guidelines for selecting test data. We can divide the input space into multi-dimensional subdomains such that the software responds to all the points within the same subdomain in a similar way. Both experience and intuitive reasoning would suggest that a significant fraction of faults would tend to occur at the boundary values. Also since the program behavior for the different points within such a subdomain is likely to be strongly correlated, testing for just a single internal point in the subdomain may be adequate in many cases. Testing for boolean conditions is discussed in [13, 14]. Automated test generation has been considered by some researchers [14, 15]. The problem of reducing the number of tests by limiting the total number of combinations to be considered by using orthogonal arrays is given in [16, 17]. This approach, although not considered here, can be used in conjunction with the scheme proposed here. One major difference between the two approaches is that antirandom testing will test for all input interactions provided sufficient test vectors are applied. It is thus applicable for ultra-high reliability software also.

There can be two possible test scenarios. For large systems, testing would have to be terminated long before it has exhausted all combinations. In this case antirandom testing would attempt to probe well distributed points, resulting in higher defect finding capability. There may be some cases (e.g. unit testing) where exhaustive testing in some sense may be possible (perhaps in terms of equivalence partitioning). In such cases, antirandom testing is likely to find defects sooner, thus reducing the overall test and debugging time.

The next section introduces and explains the basic concepts used. It also considers the problem of generating binary antirandom test sequences. In the third section, we consider automated testing of software. A checkpoint encoding scheme is introduced that reduces the general problem of testing software to generation of antirandom test sequences. Further experimental and theoretical work needed is identified in the second and the third section as well as in the concluding section.

2 Binary Antirandom Sequences

Here we start with formal definitions of the terms used and then examine construction of antirandom sequences. We assume that the input variables are all binary.

Definition: Antirandom test sequence (ATS) is a test sequence such that a test t_i is chosen such that it satisfies some criterion with respect to all tests t_0, t_1, \dots, t_{i-1} applied before.

Definition: Distance is a measure of how different two vectors t_i and t_j are. Here we use two measures of distance defined below.

Definition: Hamming Distance (HD) [18] is the number of bits in which two binary vectors differ. It is not defined for vectors containing continuous values.

Definition: Cartesian Distance (CD) between two vectors, $A = \{a_N, a_{N-1}, \dots, a_1, a_0\}$ and $B = \{b_N, b_{N-1}, \dots, b_1, a_0\}$ is given by:

$$CD(A, B) = \sqrt{(a_N - b_N)^2 + (a_{N-1} - b_{N-1})^2 + \dots + (a_0 - b_0)^2} \quad (1)$$

If all the variables in the two vectors are binary, then equation 1 can be written as:

$$\begin{aligned} CD(A, B) &= \sqrt{|a_N - b_N| + |a_{N-1} - b_{N-1}| + \dots + |a_0 - b_0|} \\ &= \sqrt{HD(A, B)} \end{aligned} \quad (2)$$

Example 1: Consider a pair of vectors:

A = {0000}

B = {1010}

Then HD(A,B) = 2 and CD(A,B) = $\sqrt{2}$ □

Definition: Total Hamming Distance (THD) for any vector is the sum of its Hamming distances with respect to all previous vectors.

Definition: Total Cartesian Distance (TCD) for any vector is the sum of its Cartesian distances with respect to all previous vectors.

Definition: Maximal Distance Antirandom Test Sequence (MDATS) is a test sequence such that each test t_i is chosen to make the total distance between t_i and each of $t_0, t_1 \dots t_{i-1}$ maximum, i.e.

$$TD(t_i) = \sum_{j=0}^{i-1} D(t_i, t_j) \quad (3)$$

is maximum for all possible choices of t_i . We will use Hamming distance and Cartesian distance to construct MHDATSs and MCDATSs.

Example 2: Consider the partial binary test sequence $t_0 = \{0,0,0,0\}$ and $t_1 = \{1,1,1,1\}$. To construct a 3-test MHDATS, t_2 can be any one of the remaining 14 vectors. Both sequences are valid MHDATSs:

$t_0 = 0\ 0\ 0\ 0$	$t_0 = 0\ 0\ 0\ 0$
$t_1 = 1\ 1\ 1\ 1$	$t_1 = 1\ 1\ 1\ 1$
$t_2 = 0\ 0\ 0\ 1$	$t_2 = 0\ 0\ 1\ 1$

For the first sequence $THD(t_2) = 1 + 3 = 4$. For the second sequence also $THD(t_2) = 2 + 2 = 4$.

However, the first sequence is not a MCDATS like the second. For the first sequence $TCD(t_2) = \sqrt{1} + \sqrt{3} = 2.732$, and for the second $TCD(t_2) = \sqrt{2} + \sqrt{2} = 2.828$. □

For black-box testing, we have no structural information available about the actual implementation. Using maximal distance criterion, every time we attempt to find a test vector as different as possible from all previously applied vectors. The antirandom testing scheme thus attempts to keep testing as efficient as possible. In this approach we are using the hypothesis that if two input vectors have only a small distance between them then the sets of faults encountered by the two is likely to have a number of faults in common. Conversely, if the distance between two vectors is large, then the set of faults detected by one is likely to contain only a few of the faults detected by the other.

If testing is less than exhaustive, then MDAT (maximum distance antirandom testing) is likely to be more efficient than either random or pseudorandom testing. This is likely to be the case for all practical systems. Even when exhaustive testing is feasible, MDAT is likely to detect the presence of bugs earlier.

When some structural information is available, it is possible to enhance MDAT using educated guesses, just as random testing can be enhanced by using weighted random testing. However, only black-box testing is considered here. A simple but computationally expensive procedure may be specified in this way.

Procedure 1. Construction of a MHDATS (MCDATS):

Step 1. For each of N input variables, assign an arbitrarily chosen value to obtain the first test vector. As discussed below this does not result in any loss of generality.

Step 2. To obtain each new vector, evaluate the THD (TCD) for each of the remaining combinations with respect to the combinations already chosen and choose one that gives maximal distance. Add it to the set of selected vectors.

Step 3. Repeat step 2 until all 2^N combinations have been used. □

This procedure uses exhaustive search. As we will see later, the computational complexity can be greatly reduced.

To illustrate the process of generating MDATS, we consider in detail the generation of a complete sequence for three binary variables.

Example 3: For a system, the inputs $\{x,y,z\}$ can be either 0 or 1. We will illustrate the generation of MHDATS using a cube with each node representing one input combination.

Let us start with the input $\{0,0,0\}$. This does not result in any loss of generality. As we will see later, the polarity of any variable can be inverted. The next vector t_1 of the MHDATS is obviously $\{1,1,1\}$ with $THD(t_1) = 3$. At this point, the situation is shown in Fig. 1a, where the input combinations already chosen are marked.

As can be visually seen, a symmetrical situation exists now. Any vector chosen would have HD = 1 from one of the past chosen vectors and HD = 2 from the others. If we allow the variables to be reordered, then without any loss of generality we have the following choices. $t_0 = \{0,0,0\}$

$t_1 = \{1,1,1\}$; THD = 3

$t_2 = \{0,1,0\}$; THD = 3 or $t_2 = \{1,0,1\}$; THD = 3

Let us consider the first choice. After $t_2 = \{0,1,0\}$, the clear choice for t_3 is $\{1,0,1\}$ at the opposite corner of the cube. The situation now is shown in Fig. 1b.

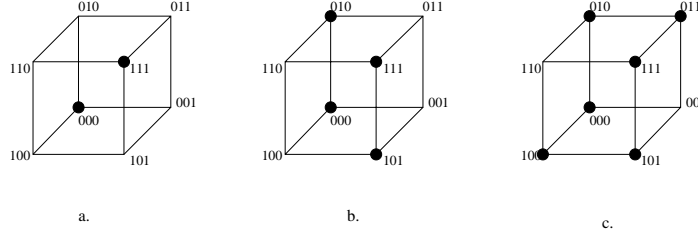


Figure 1: Construction of 3-bit MHDATS

Again a symmetrical situation exists. Any one of the remaining vectors have the same relationship with the set of vectors already chosen. Let us pick $\{1,0,0\}$ as t_4 . The next vector t_5 then has to be $\{0,1,1\}$ at the opposite corner of the cube. We can again choose any one of two remaining vectors as shown in Fig. 1c. Let us choose $t_6 = \{1,1,0\}$ which leaves $t_7 = \{0,0,1\}$. The complete MHDTS obtained here is given as sequence 1 in Table 1. \square

Table 1: 3-bit MHDTS (Example 3)

Test	xyz	THD	TCD
t_0	0 0 0		
t_1	1 1 1	3	1.7320
t_2	0 1 0	3	2.4142
t_3	1 0 1	6	4.146
t_4	1 0 0	6	4.8284
t_5	0 1 1	9	6.5604
t_6	1 1 0	9	7.2426
t_7	0 0 1	12	8.9746

In this example, it is easy to see that with our chosen vectors for t_0, t_1 and the two choices for t_2 , we could have constructed 16 distinct MHDATSs using all of the later choices available. We can verify that all of these are also MCDATSs.

A large number of experiments with construction of MHDATSs and MCDATSs have been done. Based on these, the following results can be stated.

Definition: If a sequence B is obtained by reordering the variables of sequence A, then B is a **variable-order-variant (VOV)** of A.

Theorem 1: If a sequence B is variable-order-variant of a MHDATS (MCDATS) A, then B is also a MHDATS (MCDATS).

The theorem follows from the fact that Hamming or Cartesian distance is independent of how the variables are ordered.

Example 4: The sequence A below is both a MHDATS and a MCDATS. The sequence B constructed by switching the second and the third columns is also both MHDATS and MCDATS, as can be verified calculating THD and TCD values for all options for t_1, t_2 and t_3 .

Sequence A	Sequence B
$t_0 = 0\ 0\ 0\ 0$	$0\ 0\ 0\ 0$
$t_1 = 1\ 1\ 1\ 1$	$1\ 1\ 1\ 1$
$t_2 = 1\ 0\ 1\ 0$	$1\ 1\ 0\ 0$
$t_3 = 0\ 1\ 0\ 1$	$0\ 0\ 1\ 1$

\square

Definition: If a binary sequence B is obtained by changing the polarity (i.e. inverting all the values) of one or more variables of a sequence A, then B is termed a **polarity-variant** of A.

Theorem 2: If a sequence B is a polarity-variant of a MHDATS (MCDATS) A, then B is also MHDATS (MCDATS).

The theorem follows from the fact that for a pair of vectors the distance remains the same, if the same set of variables in both are complemented.

Example 5: The sequence B below is obtained by complementing the second and the fourth variables. It can be seen that both are MHDATSs and MCDATSs. Sequence A Sequence B

$t_0 = 0\ 0\ 0\ 0$	$0\ 1\ 0\ 1$
$t_1 = 1\ 1\ 1\ 1$	$1\ 0\ 1\ 0$
$t_2 = 1\ 0\ 1\ 0$	$1\ 1\ 1\ 1$
$t_3 = 0\ 1\ 0\ 1$	$0\ 0\ 0\ 0$

□

Theorem 3: A MHDATS (MCDATS) will always contain complementary pair of vectors, i.e. t_{2k} will always be followed by t_{2k+1} which is complementary for all bits in t_{2k} where $k = 1, 2, \dots$

Proof: The first two vectors t_0 and t_1 will always be a complementary pair. Below we will show that if the vectors $t_0, t_1, \dots, t_{2k-1}$ are complementary pairs, then the vectors t_{2k} and t_{2k+1} will also be a complementary pair. Thus t_2 and t_3 , hence t_4 and t_5 , etc. will all be complementary pairs.

Let us assume a MHDATS contains an even number of vectors, $t_0, t_1 \dots t_{2k-2}, t_{2k-1}$, such that a vector with an odd subscript is a complement of the preceding vector. Let us now assume that a vector t_{2k} is found such that its total Hamming (Cartesian) distance from all of t_0, \dots, t_{2k-1} is maximal.

If we change the polarity of all the variables, we get the sequence $\bar{t}_0, \bar{t}_1, \dots, \bar{t}_{2k-2}, \bar{t}_{2k-1}, \bar{t}_{2k}$, where \bar{t}_i indicates complement of t_i . By Theorem 2, this is also a MHDATS (MCDATS). Now since $\bar{t}_1 = t_0, \bar{t}_0 = t_1$ etc., the partial sequence $\bar{t}_0, \bar{t}_1 \dots \bar{t}_{2k-2}, \bar{t}_{2k-1}$ contains the same vectors as $t_0, t_1 \dots t_{2k-2}, t_{2k-1}$ but in a different order. Thus \bar{t}_{2k} has the same total Hamming (Cartesian) distance as t_{2k} with respect to $t_0, t_1 \dots t_{2k-2}, t_{2k-1}$.

Now after having constructed the MHDATS (MCDATS) $t_0, t_1, \dots, t_{2k-2}, t_{2k-1}, t_{2k}$, the addition of $t_{2k+1} = \bar{t}_{2k}$ will keep it a MHDATS (MCDATS) because:

1. t_{2k} provides the maximal total HD (CD) with respect to $\{t_0, t_1 \dots t_{2k-2}, t_{2k-1}\}$.
2. \bar{t}_{2k} also provides the maximal HD (CD) with respect to t_{2k} . QED.

□

So far we have assumed the construction of MHDATSs (MCDATS) using exhaustive evaluation of all the remaining vectors. Thus for a vector containing N variables, the computation will require evaluation of $\{1 + (2^N - 1) + (2^N - 2) + \dots + 4 + 3 + 2 + 0\}$ distances to obtain a complete sequence. Using the above property, the number of evaluations is reduced to $\{1 + 0 + (2^N - 2) + 0 + (2^N - 4) + \dots + 4 + 0 + 0 + 0\}$ distances. That may still represent a prohibitive amount of computation in situations when N is large. We will next consider incremental construction of complete MHDATSs and MCDATSs.

Procedure 2. Expansion of MHDATS (MCDATS):

Step 1. Start with a complete MHDATS of N variables, $X_{N-1}, X_{N-2}, \dots, X_1, X_0$.

Step 2. For each vector $t_i, i = 0, 1, \dots, (2^N - 1)$, add an additional bit corresponding to an added variable X_N , such that t_i has the maximum total HD (CD) with respect to all previous vectors. □

We have extensively experimented with this procedure. We have two major observations.

1. It is always possible to extend a MHDATS (MCDATS) by adding one more variable using Procedure 2. A formal proof for this is being sought.
2. The column added is a function of columns already included. We are attempting to identify this function such that it does not vary for different values of N and can be used in a simple way.

To illustrate additional aspects of constructing antirandom sequences, let us consider the following example.

Example 6: Here let us attempt to construct a 4-variable antirandom test sequence. We will use Procedure 2 to extend the 3-variable sequence obtained in Example 3 (Table 1) by adding an extra column. The complete 4-variable sequence will have 16 vectors as opposed to the 3-variable sequence we have with 8 vectors. We will do

the construction in two parts. First, we will augment the 3-variable/8 vector sequence into a 4-variable/8 vector sequence. We will then construct the rest of the 8 vectors.

a. Let us append a fourth variable w to the lefthand side. If we start by adding a '0' to t_0 , we have to add a '1' to t_1 , then for t_2 we have two choices.

$$\begin{array}{c} \text{w x y z} \\ t_0 = 0 0 0 0 \\ t_1 = 1 1 1 1 \end{array}$$

$$t_2 = 0 0 1 0 \text{ or } t_2 = 1 0 1 0$$

With the first choice we have $\text{THD} = 1 + 3 = 4$ and $\text{TCD} = \sqrt{1} + \sqrt{3} = 2.7320$. The second choice gives us $\text{THD} = 2 + 2 = 4$ again, but $\text{TCD} = \sqrt{2} + \sqrt{2} = 2.8284$. Since the second choice would generate a sequence that is a MCDATS as well as a MHDATS, let us use it for t_2 . The vector t_3 then is $\{0 1 0 1\}$.

Using Cartesian distance as a criterion, we can construct the next four vectors similarly. the partial sequence then is given below. We have used up the complete 3 bit sequence given in Table 1. The column corresponding to the added variable is shown in bold.

$$\begin{array}{l} t_0 \quad \mathbf{0} 0 0 0 \\ t_1 \quad \mathbf{1} 1 1 1 \\ t_2 \quad \mathbf{1} 0 1 0 \\ t_3 \quad \mathbf{0} 1 0 1 \\ t_4 \quad \mathbf{1} 1 0 0 \\ t_5 \quad \mathbf{0} 0 1 1 \\ t_6 \quad \mathbf{0} 1 1 0 \\ t_6 \quad \mathbf{1} 0 0 1 \end{array}$$

b. Now a symmetrical situation exists. We can use any of the remaining vectors for t_8 ; however, it would be useful if certain patterns are repeated for the next eight vectors.

Our extensive experimentation suggests that we can construct the rest of the sequence by taking the above partial sequence and simply changing the polarity of leftmost variable w. The last part of the sequence then is:

$$\begin{array}{l} t_8 \quad \mathbf{1} 0 0 0 \\ t_9 \quad \mathbf{0} 1 1 1 \\ t_{10} \quad \mathbf{0} 0 1 0 \\ t_{11} \quad \mathbf{1} 1 0 1 \\ t_{12} \quad \mathbf{0} 1 0 0 \\ t_{13} \quad \mathbf{1} 0 1 1 \\ t_{14} \quad \mathbf{1} 1 1 0 \\ t_{15} \quad \mathbf{0} 0 0 1 \end{array}$$

The calculation of TCD and THD proves that the 16-vector sequence is indeed a MCDATS and a MHDATS. \square

The Example 6 above represents one of many such construction experiments. The observations are summarized below.

1. A MCDATS is always a MHDATS but not vice versa. For the same number of bits there are fewer valid MCDATSs than MHDATSs.
2. When constructing an N-variable sequence, a symmetrical situation exists, when the first $(2^N/2)$ vectors have been obtained, i.e. any one of the remaining vectors can be chosen as the next vector.
3. For N-variables, the first $(2^N/2)$ vectors are such that if the polarity of a variable is reversed, then the resulting vectors $(2^N/2)$ are distinct from the original set of $(2^N/2)$ vectors.
4. For $2^N/2$ vectors of an N-bit MCDATS (MHDATS), we can obtain the rest of the $2^N/2$ vectors by changing the polarity of one of the variables.

The formal proofs for these are being sought. The above observations suggest an incremental procedure to construct an N-variable MHDATS (MCDATS).

Procedure 3. Expansion and Unfolding of a MHDATS (MCDATS):

Step 0. Start with a complete (N-1) variable MHDATS (MCDATS) with 2^{N-1} vectors.

Step 1. Expand by adding a variable using Procedure 2. We now have the first $(2^N/2)$ vectors needed.

Step 2. Complement one of the columns and append the resulting vectors to first set of vectors obtained in Step 1. Here, it would be convenient to complement the variable added in Step 1. \square

The procedure requires $\frac{2^N}{2} \cdot \frac{1}{2} = 2^{N-2}$ bits to be evaluated.

In many cases we would like to obtain an antirandom sequence involving a large number of variables. Instead of incrementally expanding and unfolding, can we concatenate m N -bit antirandom sequences to obtain a $m \times N$ bit antirandom sequence with 2^N vectors? Can we obtain the rest of $(2^{m \cdot N} - 2 \cdot N)$ vectors in an algorithmic manner to create the complete sequence? Consider the following example.

Example 7: Let us concatenate two copies of the 3-bit antirandom sequence to obtain the following 6 bit sequence.

t_0	0 0 0 0 0 0
t_1	1 1 1 1 1 1
t_2	0 1 0 0 1 0
t_3	1 0 1 1 0 1
t_4	1 0 0 1 0 0
t_5	0 1 1 0 1 1
t_6	1 1 0 1 1 0
t_7	0 0 1 0 0 1

An evaluation shows this to be a MHDATS but not a MCDATS. We can easily see that for a MCDATS, the vector t_2 should have an equal number of ones and zeroes. \square

Our experiments suggest that such concatenation should retain the maximal HD property. We need to obtain a formal proof for this. We also need to obtain an algorithmic method to obtain wide MCDATSs and to obtain the rest of the $(2^{m \cdot N} - 2^N)$ vectors. It is our conjecture that it should be possible by using polarity variants of the component arrays.

In some applications it may be possible to order and partition the set of variables into group of size p , such that effect of variables in one group is disjoint (or nearly disjoint) to the variables in all other groups. In this case, all we need to do is to apply an p -variable antirandom sequence to each group in parallel. This can considerably reduce the test application time.

3 Checkpoint Encoding

In a general case, the inputs can be numbers, alphanumeric characters as well as data structures composed using them. In such cases also we would like to maximize the effectiveness of testing. It is possible for one to define ‘distance’ and use them for constructing antirandom sequences in such cases also.

Example 8: Let us consider two real variables x and y which can range from 0 to 1. The following then is a MCDATS.

0	0
1	1
0.5	0.5

\square

However, defining ‘distance’ can be difficult for data structures. Also for a program, the input variables can be of different types and ranges, which will make construction of antirandom sequences extremely hard. We here propose an encoding approach which will convert the problem to constructing binary antirandom sequences. The approach is based on **domain** and **partition analysis** and the concepts of **equivalence partitioning**, **revealing subdomains** [11] and **homogeneous subdomains** [19]. The technique partially encodes an input into binary, such that sample points desired can be obtained by automatic translation.

These sample points, termed **checkpoints** here, are strategically selected such that they are likely to span most types of variations in the program behavior with respect to each input. To illustrate the approach let us consider this simple example. For convenience, we use a square bracket (“[” or “]”) to indicate inclusion of the

endpoint and a parenthesis (“(” or “)”) to indicate exclusion.

For testing, it is important to test for illegal input values because the program must respond correctly to those inputs, as we see in the following example. The range of illegal values should be regarded as one or more additional equivalent partitions.

Example 9: Consider a continuous input variable x which can range from x_{min} to x_{max} , with the end values included. A value outside of $[x_{min}, x_{max}]$ is illegal (Figure 2). Let us call all the values of x such that $x_{min} < x < x_{max}$ **internal points**. Let us assume that the program under test either works correctly or incorrectly for all internal points. We can thus use the following **checkpoint encoding** scheme as shown in Table 2.

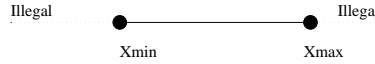


Figure 2: Normal and illegal ranges

Table 2: Encoding scheme for Example 9

Encoded Binary	Value of x
0 0	randomly chosen value in the illegal range
0 1	x_{min}
1 0	x_{max}
1 1	a randomly chosen internal point

□

In general encoding can take several considerations into account based on specifications and the fault hypothesis used.

1. Some fraction of all input values applied should be illegal.
2. Many defects may be associated with boundary values. They may not be detected when “typical” values are used for testing. Some defects may require testing points closest to the boundary on both sides.
3. A legal range of values can often be partitioned into subdomains, such that each subdomain exercises a somewhat different (but not necessarily disjoint) part of the functionality. Thus one can use a fault hypothesis that exercising only one such subdomain, defects corresponding to other subdomains may not be triggered.
4. With black-box testing, specifications may not reveal the reasons which would suggest a subdomain should have been further partitioned. Thus there may be a need to sample a few randomly chosen points within a subdomain. In practice, some information about the software structure can be useful for proper identification of subdomains.
5. If some inputs are more critical, there may be a reason to use a larger number of samples for that input. This can be done by using more bits to encode that input.

Let us now see an example that illustrates some of the above considerations.

Example 10: Consider a continuous input variable x which is valid in the range $[-a, b]$. Let us assume the program may respond differently in subdomains $[-a, 0)$ and $(0, b]$. In addition, we also wish to sample the function with a couple of values each within both $[-a, 0)$ and $(0, b]$.

Here we can use the checkpoint encoding scheme shown in Table 3. It attempts to ensure that the special cases (boundary and illegal values) are generated early and the adjacent internal points are also adjacent in the Hamming distance sense.

□

After we have encoded all the inputs using binary variables, the problem is simply reduced to that of generating binary antirandom sequences. If the input variable I_j $j = 0, 1, \dots, M$ where M is the total number of inputs, is encoded using C_j bits, the binary vectors would be $\sum_{j=0}^m C_j$ bits wide. Notice that the tests generated by this

Table 3: Encoding scheme for Example 10

Encoded Binary	Value of x
1 0 1	randomly chosen in the illegal range
0 0 0	$-a$
1 0 0	randomly chosen in $(-a, -a/2]$ range
1 1 0	randomly chosen in $(-a/2, 0)$ range
0 1 0	0
0 1 1	randomly chosen in $(0, b/2)$ range
0 1 1	randomly chosen in $[b/2, b)$ range
1 1 1	b

scheme will include common combinations, likely to be encountered most often during operational use as well as special combinations which have higher error revealing capability.

The proposed scheme, as illustrated in Figure 3, combines the benefit of checkpoint encoding with antirandom testing. Once the checkpoint encoding definition (CED) has been obtained, the tests can be generated automatically. The antirandom vectors generated are translated using the CED. When needed, the random value generator will generate a random value of an input, which may be a single number or a data structure. The advantage of randomization is that when a subdomain for an input variable is encountered several times, each time a different value is likely to be generated. This would require that the seed values be altered during each access.

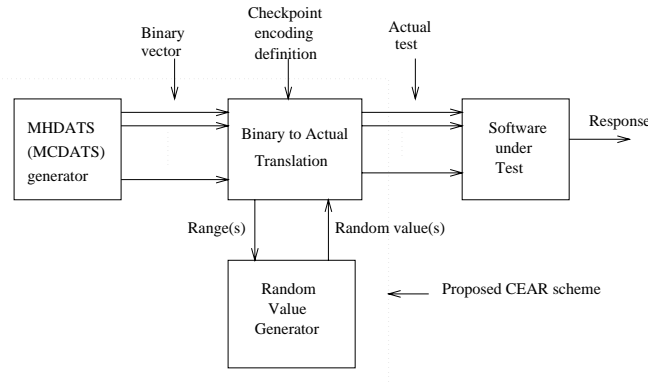


Figure 3: Proposed CEAR scheme with encoded antirandom vectors.

The computational requirements of the blocks included in the Checkpoint Encoded Antirandom testing (CEAR) scheme are very light. Thus it may not be necessary to record the set of tests before they are applied. The CEAR and the software under test may run together, a test being used right after it is generated. The CEAR scheme may be implemented to make the test-suite generated either repeatable or nonrepeatable. The first option would be suitable if a fixed regression testing test-suite is desired.

We must discuss a very important tradeoff. Many defects can be associated with boundary conditions. Such a condition can be subdomain boundary for a single input variable or a combination of such boundaries for several input variables. On the other hand, in normal operation, much of the time only internal values will occur. Thus the field failure rate may depend more on the faults associated with the values which are not boundary conditions.

In terms of the checkpoint encoding scheme, the question can be stated in this way. For each variable that is part of the input space, how many checkpoints should represent ‘common’ values, and how many should be boundary values? We perhaps require a considerable amount of experimental data before we can obtain definite guidelines for answering this question. If we can be certain that there are no unidentified subdomain boundaries (that will however require examining the software implementation), we could assume that each internal subdomain is a revealing subdomain, and thus it needs to be sampled only once. For very small software systems, it may be possible to exhaustively test for all boundary values combinations along with a reasonable number of common combinations. However, for a general system, guidelines for making optimal choices cannot be offered at this time.

Example 11: As a more complete example of checkpoint encoding, let us consider the FIND program by Hoare (Figure 4). Testing of its FORTRAN version was examined by DeMillo et al. [8]. It takes as an input an integer array A with size $N \geq 1$ and an array index F, $1 \leq F \leq N$. The program rearranges the array such that all elements to the left of A(F) have values no larger than A(F), and all elements to the right are no smaller.

The subdomains are identified below; the special cases are marked with an asterisk.

Array size: 1, 2 and greater than two

Element values:

- all positive,
- all 0,
- all negative,
- illegal (containing non-integers),
- mixed.

F points to:

- first element,
- a middle element,
- last element
- outside (illegal)

Array status:

1. elements randomly ordered
2. elements already ordered as needed
3. elements in order reverse of case 2
4. elements all equal

In order to encode the checkpoints efficiently, we may want to separate special cases and consider them separately. For example When the array size is 1, the choices for F and array status do not make sense. Similarly the case when all elements are 0 may also be considered separately. Separating such cases, we may use the encoding scheme suggested in Table 4. It can be called a "field encoding scheme" because each field of a few bits is encoded separately.

Table 4: Encoding scheme for Example 11

Field	Bits	Value	Significance
Array Size	b0	0	2
		1	> 2
Element values	b3,b2,b1	000	all positive
		001	all negative
		010	illegal
		rest	mixed
F points to	b5,b4	00	first element
		01	a middle element
		10	last element
		11	illegal
Array status	b7,b6	00	randomly ordered
		01	already ordered
		10	reverse ordered
		11	all equal

We have thus used eight variables for encoding. A partial nine vector 8-bit antirandom sequence is shown in Table 5. It should be noticed that the first 8 rows are formed by concatenation of two partial 4-bit antirandom sequences. It is possible to use some multibit combinations across several fields to specify special cases.

Notice that all the individual choices have occurred in the first eight vectors, but to apply all possible combinations allowed by the encoding scheme, 2^8 vectors would be needed. \square

It can be seen that the resulting combinations are consistent with the findings and recommendations of DeMillo et al. [8]. They suggest use of illegal as well as special values in addition to the usual values. If additional testing time is available, the coding scheme above should be modified to allow a larger number of normal combinations.

Table 5: A partial 8-bit sequence

b7	b6	b5	b4	b3	b2	b1	b0
0	0	0	0	0	0	0	0
1	1	1	1	1	1	1	1
1	0	1	0	1	0	1	0
0	1	0	1	0	1	0	1
1	1	0	0	1	1	0	0
0	0	1	1	0	1	1	0
1	0	0	1	0	0	1	1
0	1	1	0	0	1	1	0
1	0	0	1	1	0	0	1
0	1	1	1	1	0	0	0

```

SUBROUTINE FIND(A,N,F)
C
C FORTRAN VERSION OF HOARE'S FIND
C PROGRAM (DIRECT TRANSLATION OF
C THE ALGOL 60 PROGRAM FOUND IN
C HOARE'S "PROOF OF FIND" ARTICLE
C IN CACM 1971).
  INTEGER A(N),N,F
  INTEGER M,NS,R,I,J,W
  M=1
  NS=N
10 IF(M.GE.NS) GOTO 80
  R=A(F)
  I=M
  J=NS
20 IF(I.GT.J) GOTO 60
30 IF(A(I).GE.R) GOTO 40
  I=I+1
  GOTO 30
40 IF(R.GE.A(J)) GOTO 50
  J=J-1
  GOTO 40
50 IF(I.GT.J) GOTO 20
C
C COULD HAVE CODED GO TO 60 DIRECTLY
C -DIDN'T BECAUSE THIS REDUNDANCY
C IS PRESENT IN HOARE'S ALGOL
C PROGRAM DUE TO THE SEMANTICS OF
C THE WHILE STATEMENT.
C
  W=A(I)
  A(I)=A(J)
  A(J)=W
  I=I+1
  J=J-1
  GOTO 20
60 IF(F.GT.J) GOTO 70
  NS=J
  GOTO 10
70 IF(I.GT.F) GOTO 80
  M=I

```

```
GOTO 10
80 RETURN
END
```

Figure 4. FIND program from [8].

4 Conclusion and Future Work

We have presented a black-box approach that attempts to maximize the test effectiveness by keeping tests as different as possible from each other. The scheme provides formalization of a concept that is intuitively attractive. Unlike coding theory or random/pseudorandom testing, this is a new approach that requires further theoretical and experimental investigations. Some of the theoretical challenges are identified in the paper. We need to experimentally evaluate and refine the techniques for well known benchmark programs for comparison and characterization as well as for larger systems. The proposed technique needs to be compared with existing white-box and random testing approaches. Techniques that allow application of this approach to large software systems need to be developed and studied. The checkpoint encoding scheme proposed here converts the test generation requirement to a binary problem. It also has the advantage of explicitly enumerating the checkpoints which is vital to keep testing efficient.

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