
Simulating a General Purpose Mobile Computing Environment*

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July 9, 1998

Technical Report CS-98-110

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Abstract

This paper reports the modeling and implementation of a general purpose mobile computing environment. The design is modular and it incorporates all intrinsic operations of mobile computing, using the working of a mobile telephony system as a basis. The SES/Workbench has been used as the primary tool for implementation; one of the purposes was to investigate the usefulness of the tool to simulate a scalable and flexible mobile computing testbed. Models and submodels have been described and the correctness and the scalability of the design have been demonstrated by presenting experimental results of simulating a *fixed channel assignment* scheme. computing environment. The simulation has been designed in a modular fashion incorporating all the intrinsic operations of mobile computing, taking the working of a mobile telephony system as a basis. To demonstrate the correctness, scalability and flexibility of our simulation model we present results for a *fixed channel assignment* scheme.

1 Introduction

Mobile computing is a distributed computing paradigm where users are free to move about without losing the ability to connect to a network. This poses several interesting problems which are different from those in a static network. Routing a connection to a mobile host involves determining

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its location. Communication takes place over wireless channels. Wireless channels are characterized by low bandwidths and are less reliable than their static counterparts due to their susceptibility to interference and fading. There is a need to have channel allocation schemes that mitigate these effects[Jor95]. Mobile hosts are relatively underprivileged compared to nodes on a static network in terms of computational power and network connectivity. Innovations are required to reduce the computational requirements by shifting some of the computational load to static hosts. Caching schemes to lower the reliance on the network connectivity also need to be looked into. Hence there is a need for a comprehensive testbed to study the effects of mobility and determine the effectiveness of all such proposed schemes. In this project we have developed a modular simulation environment for mobile computing. To develop the simulation we used SES/Workbench, which is an event driven simulation tool [Sci94a] [Sci94b]. Our simulation model is flexible, allowing us to easily simulate proposed solutions to various mobile computing problems. It is scalable in that it can simulate networks of different sizes with different parameters.

The rest of this paper is organized as follows. In the next section we explain the system model used in our mobile computing environment. Section 3 explains the high level protocol that we used for connection setup and mobility. Section 4 gives a brief explanation of the various modules that make up our simulation. Each of these have a corresponding implementation in SES/Workbench. Section 6 presents an example simulation to illustrate the current status of the simulator and finally section 7 gives conclusions and the enhancements that still need to be made to the simulator.

2 System Model

A mobile system consists of the following entities

- **Cell** A geographical area consisting of a Mobile Support Station(MSS) and a dynamic set of Mobile Hosts(MHs) communicating with the mobile support station. A cell is characterized by it's geographical area, the MSS that caters to the MHs in this area and a set of wireless channels that are used by the MSS to communicate with these MHs. Two neighboring cells

may not be able to use the same set of wireless frequencies as they might interfere with communications in an adjoining cell.

- **Mobile Host (MH)** A user's portable device e.g. a laptop computer, a cell phone, that enables him to connect to the network.
- **Mobile Support Station (MSS)** A node on the static network which provides an interface to the static network for each MH in it's cell. Communication requests by the MH are routed to the MSS via a wireless link. The MSS is responsible for routing the request on the wired network and sending the result of the request back to the MH.

3 Protocol

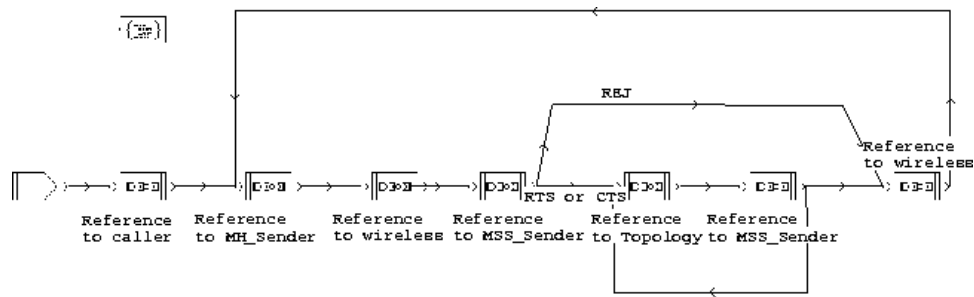


Figure 1: Call Module

3.1 Call Setup

- MH initiating a call sends an RTS(Request To Send) to its MSS
- MSS allocates a channel to the MH (*Channel Allocation*) if its available, otherwise sends back an REJ(REJECT).
- After allocating a channel, MSS locates the present position of the callee MH (*Location Management*) and sends an RTS to the callee's MSS.

- The callee's MSS allocates a channel for the call, if available. If not, it routes an REJ back to the caller MH via the caller's MSS.
- After allocating a channel, the callee MSS sends an RTS to the callee MH.
- The callee MH responds with a CTS (Clear To Send) and routes it back to the caller MH, retracing the path of the RTS.
- The connection is setup once the CTS reaches the caller MH. Both parties can now send data.
- Either party can stop the communication by sending a disconnect message and channels are freed at both ends.

3.2 Handoff / Call maintenance

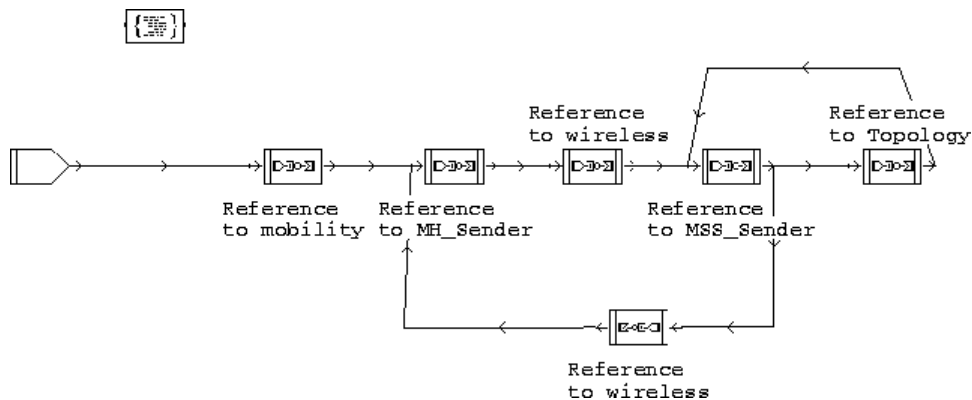


Figure 2: Mobility Module

- MH moving from one MSS to the other (MH_{mover} moving from MSS_{source} to MSS_{dest}) sends a *greeting* message to MSS_{dest} .
- If a call is in progress, MSS_{dest} allocates a channel for this call, if no channel is available, a disconnect message is sent to both caller and callee MHs.
- MSS_{dest} sends a *deregister* message to MSS_{source} .

- On receiving the deregister message, MSS_{source} removes MH_{mover} from it's list of MHs in its cell and sends back a *register* message to MSS_{dest} . If a call for MH_{mover} is in progress, a *call maintenance* message to the correspondent MSS.
- On receiving the register message, MSS_{dest} adds MH_{mover} to the list of MHs in its cell. This concludes the Handoff process.
- On receiving a call maintenance message the correspondent MSS updates the location information for MH_{mover} from MSS_{source} to MSS_{dest} .

4 Simulation Modules

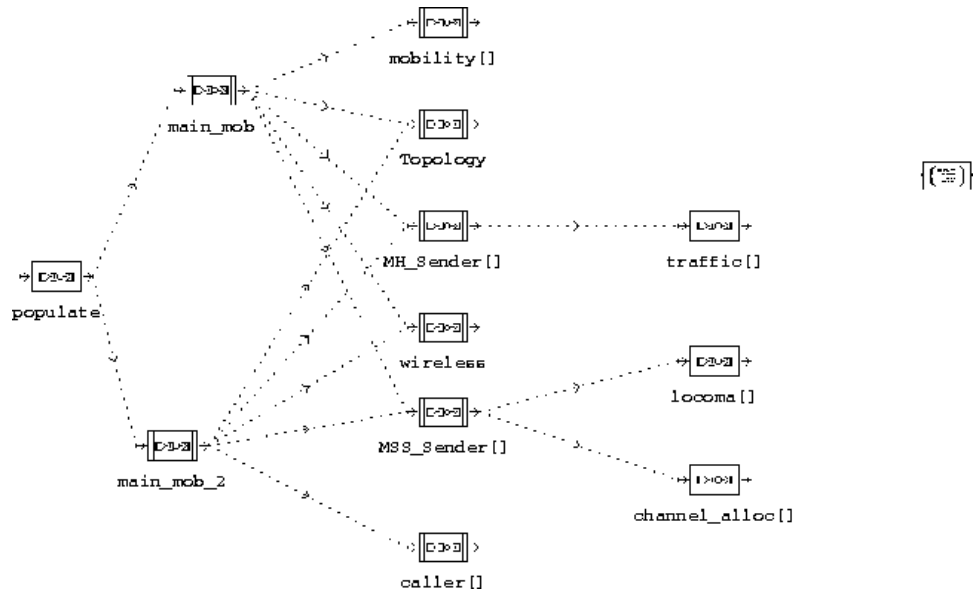


Figure 3: System Modules

The model comprises of two independent main modules - the first models mobility of users and the second refers to call generation and transmission. A population initialization module defines the initial state of the above modules.

4.1 Populate

This submodule generates the initial population of Mobile Hosts(MH) in the system. MHs belonging to various categories depending upon their mobility and calling patterns can enter the system.

4.2 Mobility Model

This submodule simulates the moves MHs make within the system. Each MH leaves its present position with a certain probability and updates its current location w.r.t. an MSS i.e. a mapping of MH to MSS is maintained.

4.3 Call Model

This is the module responsible for generating calls. It captures the calling behavior of each MH by assigning a distribution to the number of calls made by it.

4.4 Wireless

Wireless module simulates the wireless communication channel between an MH and an MSS. Messages generated by the traffic module of the MH are transmitted through the wireless channel to the MSS. Each call is given a dedicated channel at each of the two MSSs.

4.5 Mobile Service Station (MSS)

The MSS handles all communication of the MH by executing the message passing protocols described in the previous section. The MSS also has two sub-modules responsible for channel allocation and location management.

- Channel Allocation

When an MH wishes to communicate, a request is sent for channel assignment. The channel allocation module handles this request by assigning a certain free channel to the requesting MH. If a channel is not available, then the connection request by the MH is rejected.

- Location Management

After a channel has been allocated, a request is sent to determine the location of the recipient of the call. The Location Management module handles this request by running an algorithm to identify the MSS where the callee is residing and communicates this information back to the caller. these steps, the MSS routes messages to the destination through the topology.

4.6 Mobile Host (MH)

This module models the behavior of a MH in the system. Each MH is characterized by a certain calling and mobility behavior. An MH generates calls in the system according to its calling characteristic and this functionality is handled by the traffic sub-module.

- Traffic

Once a call has been established, the traffic modules at each of the correspondent MHs take over and are responsible for sending data. Characteristics of the generated traffic are determined by the type of call being executed and this information is passed down by the caller module.

4.7 Topology

This module simulates the wired network of all MSSs in the system. Messages are routed from one MSS to another, dependent on the topology.

5 SES Implementation

Each of the modules enumerated above is implemented as a submodel in SES. SES uses an object oriented approach, defining a set of nodes and transactions to model system behavior. Different nodes have different functionality e.g. a blocking node to block transactions till a condition is satisfied, a delay node to delay the transaction for a specific simulation time interval etc. Transactions follow a predefined path between nodes, executing specific code at each stage.

In our implementation, all protocol messages are modeled as transactions of different categories. The population module creates two transaction for each MH comprising the initial population. These transactions are sent to the caller and mobility modules which are dimensioned on the number of MHs as each MH has been modeled to have an independent call generation and mobility pattern. New transactions are created after inter move and inter call time delays (based on a probability distribution) by each of these modules to represent the movement (*mobility*) or call initiation (*connect*) for a MH. These transactions flow through the various modules executing the handoff and call setup phases of the protocol described above. For example, transactions of category type *mobility* are passed to the MH_model of the same dimension. The MH_model changes the type of the transaction to *greeting* and sends it to the MSS_model of the dimension corresponding to the MSS of the MH's current location.

An interesting feature of the MSS_model is the use of timeouts, even though we are assuming that messages are passed on the wired and wireless networks reliably with no errors. This feature is used due to mobility. It may so happen that when an MSS sends a message to a MH, the MH may have moved out of its current location. The MH may never see the message, thus a timeout mechanism is required to ensure that the MSS is informed of this situation and can take corrective action. This is implemented by keeping a copy of the transaction sent to the MH at a timeout delay node. If an acknowledging transaction is received before the timer expires, the transaction is sunk, otherwise the transaction flows through nodes which take corrective action according to the protocol.

5.1 Data Structures

The following data structures are maintained to model the system.

1. **MH_state**

An array of structures in which each structure element represents the state of an MH. It's size equals the number of MHs in the system. The fields of the structure are

- **MSS_id**

The id of the MSS in whose cell, the MH currently is.

- **state**

One of either callee, caller or idle.

- **MH_partner**

If the state is not idle this field has the id of the MH, the MH is communicating with.

- **MSS_partner**

Current location of MH_partner. This is set by the location management module for the caller MH.

- **established**

Is set to 1 once the call initiation phase is complete. This field is used to avoid sending unnecessary disconnect messages due to mobility when the call is still in its initiation phase.

2. **MSS_Record**

An array for each MSS in the system. Lists the ids of MHs currently local to the MSS.

3. **Packet**

A structure associated with each transaction. Has the following fields

- **MH_source**

The MH initiating the transaction.

- **MSS_source**
The MSS from whose cell the MH originated the transaction.
- **MH_dest**
The intended recipient MH of the packet.
- **MSS_dest**
MSS where MH_dest is located.

6 Example Simulation

We carried out an example simulation using our simulator modeling a fixed channel allocation scheme[JS96]. Each MSS is assumed to have a fixed number of channels which it can use to support calls originating from its cell. We made some further simplifying assumptions

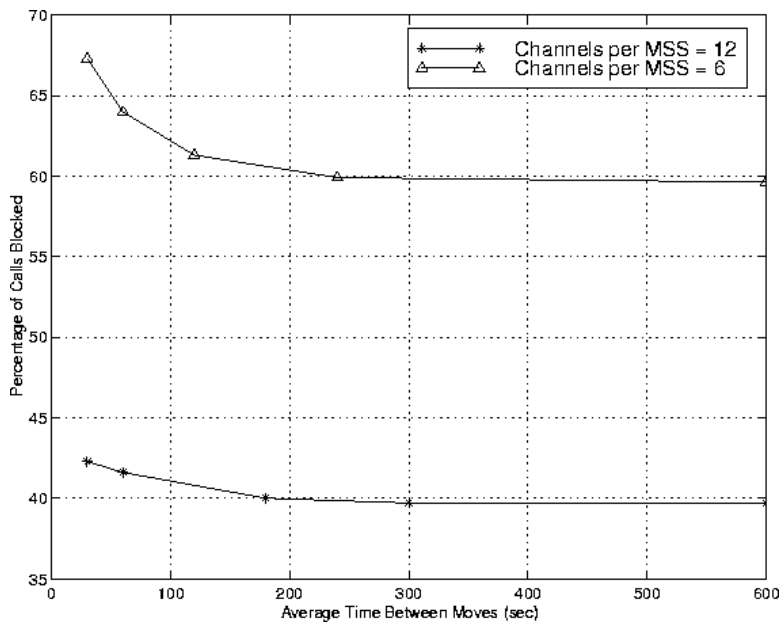


Figure 4: Mobility Vs Percentage of calls blocked

1. All users have the same mobility pattern. The time between moves for each user is exponentially distributed.

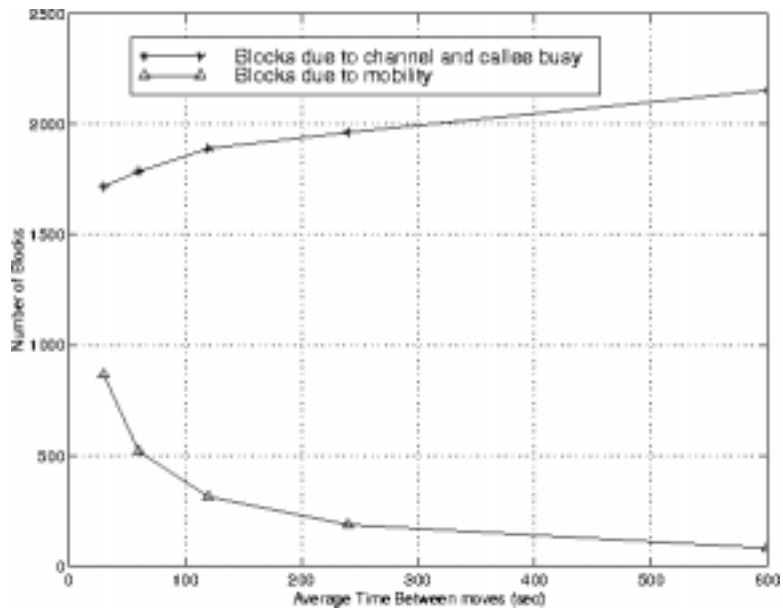


Figure 5: Comparative Distribution of Blocking

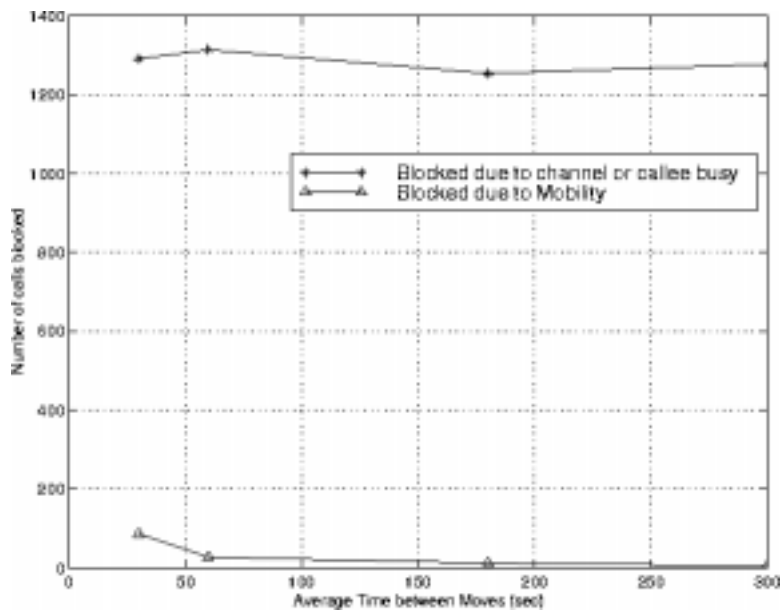


Figure 6: Comparative Distribution of Blocking

2. All users have the same calling pattern. No calls are attempted when the user is busy. The time between calls for a user is exponentially distributed from the time his state becomes idle.
3. The duration of a call is exponentially distributed.

4. There is a point to point link between any two MSSs. (This would be the case for example when all MSSs are connected to an ATM switch.)
5. Location management is perfect. There is no overhead due to location management.
6. The calls between two MHs are for voice. Voice data is modeled as a a continuous stream computed on the basis of 8 samples per second of 8 bits each. Also the intra packet end to end delay must be less than 0.5 seconds for the voice quality to be satisfactory.
7. Wireless channels are 64Kbps each.

The system parameters were taken as follows

Parameter	Distribution	Default Value
Number of MSS	-	50
Number of MHs	-	700
Number of Channels / MSS	-	12
Control message length	-	320 bits
Voice Packet length	-	16000 bits
Time between voice packets	-	0.25s
Mobility Rate	exponential	5 minutes
Calling Rate	exponential	5 minutes
Call duration	exponential	3 minutes
Channel capacity		64Kbps
Wired Links bandwidth		10 Mbps
Simulation Time		20min

6.1 Result

Calls are blocked, due to either the callee being busy, non availability of channels when initiating the call, or disconnection of calls due to a free channel not being available on a busy MH's move to

a new cell. As the mobility increases the number of calls disconnected due to mobility increases 6 whereas blocks due to other factors shows a decline. The overall call blocking rate also increases with mobility 4. The increase in blocking probability is significant over a small range of inter-move delay values. Thereafter, the effect is less pronounced as expected. Varying the number of channels from 6 to 12 decreases the call blocking probability by almost a third.

7 Future Work

The simulation system developed as part of this coursework is highly flexible and scalable. It has the potential to support different Channel Allocation, Location Management, Caching - Prefetching techniques with different mobility and calling characteristics. All this has been done in a modular fashion so that a change in one module can be done without altering any other modules.

Several features are yet to be implemented. Currently the simulator models only voice data. The topology and wireless modules are modeled as simple delays assuming dedicated channels. Call duration, call arrival rate and mobility patterns are modeled as exponential distributions. These assumptions may have to be modified depending on the model under consideration. For example, if a contention protocol is used on the wireless channel the assumption regarding dedicated channels does not hold anymore. Similarly, a Markov process may not be a realistic model for mobility as users do not have totally random movement patterns. For example, a user would move from his place of work to home and back at fixed times during the day.

In the above example simulation we saw that the blocking probability is highly dependent on the state of the callee when a caller makes a call. Callee sets should not necessarily be chosen at random, as users have a fixed set of acquaintances, whom they call most often, besides making a few random calls. Our emphasis for now was to provide a top level framework. The way the simulation is designed, there are hooks to to modify/implement features, to simulate various solutions to problems posed by mobility. This work provides us with the opportunity to compare and contrast various schemes proposed in literature and then implement some of our own by learning from the

experience.

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