Looking for something?
Trickle down from the root
Branching choices at the fork
A left if it’s less
A right if it’s more
A dead end?
What you seek
Isn’t here

Frequently asked questions from the previous class survey

- How long do records typically stay in the DNS cache?
Topics covered in this lecture

- Binary Search Trees
- B-Trees

To dwellers in a wood, almost every species of tree has its voice as well as its feature.

Thomas Hardy, Under the Greenwood Tree
Binary Search Tree (BST)

- Binary search trees use the concepts underpinning the binary search algorithm to create a **dynamic data structure**
  - The key word here is dynamic

- Unlike sorted arrays, binary search trees support the efficient **addition and removal** of elements in addition to searches
  - Making them the perfect blend of the algorithmic efficiency of binary search and the adaptability of dynamic data structures

Tree Structures

- Trees are **hierarchical** data structures
- Comprise **branching chains** of nodes
- Natural **extension** of linked lists
  - In the case of BSTs, each tree node is permitted two next pointers
    - Point to subsequent nodes in disjoint lists
Example BST

- Root Node: 8
- Internal Nodes: 3 and 23
- Leaf Nodes: 1, 5, and 11

BST Nodes

- A node contains a **value** (of a given type)
  - Plus, **up to two** pointers to lower nodes in the tree
- Nodes with at least one child?
  - Internal nodes
- Nodes without any children?
  - Leaf nodes
Tree nodes may contain other information

- Often include a pointer back to the node’s parent, for instance
  - Allows the bottom-up traversals of the tree
    - In addition to the typical top down
  - Comes in handy when we consider removing nodes

The TreeNode data structure

```
TreeNode {
  Type: value
  TreeNode: left
  TreeNode: right
  TreeNode: parent
}
```
We might also want to store auxiliary data …

- Storing and searching for individual values are useful

- However, **using these values as keys** for looking up more detailed information greatly extends the power of the data structure
  - For e.g., coffee names as the node’s values, allowing us to efficiently look up records for any coffee
    - Our auxiliary data would be a detailed record of that coffee

BST and auxiliary data

- The tree node data structure can either store this auxiliary data
  - Directly
  - Include a pointer to a composite data structure
    - Located somewhere else in memory
Binary search trees

- Start at a single root node at the top of the tree
- **Branch** into multiple paths as they descend
- Allows programs to access the binary search tree through a single pointer
  - The location of its root node

Nodes and dispersion in memory

- A search tree’s individual nodes can be **scattered** throughout memory
- Each node is only linked to its children and parents through?
  - The power and flexibility of **pointers**
The power of the binary search tree stems from how values are organized within the tree.

- For any node \( N \)
  - The value of any node in \( N \)'s left subtree is less than \( N \)'s value
    - Values in the left node and all nodes below it are less than the value of the current node
  - The value of any node in \( N \)'s right subtree is greater than \( N \)'s value
    - Values in the right node and all nodes below it are greater than the value of the current node
- The rule defines the tree's structure below that node
  - Partitions the subtree into two subsets

**Useful Rule:** The number of right-handed people is greater than the number who are left-handed.

Values in a BST are ordered by the binary search property

![Binary Search Tree Diagram](image-url)
Values in a BST are ordered by the binary search property

- This ordering of nodes might not seem like a lot of structure
  - Recall the power we got from using a similar property within binary search
- The binary search tree property is effectively keeping the data within the tree sorted with respect to its position in the tree
- As we will see, this allows us to not only efficiently find values in the tree but also efficiently add and remove nodes

Searching Binary Search Trees

- We search the BST by walking down from the root node
- At each step, we determine whether to explore the left or right subtree
  - By comparing the value at the current node with the target value
    - Left: If the target value is less than the current value
    - Right: If the target value is greater than the current value
- The search ends when either the target value is found, or it reaches a node with no children in the correct direction
  - In the latter case, we can assert that the target value is not in the tree
Searching Binary Search Trees: An analogy

- The node’s value thus serves the same function as those helpful signs in hotels
- Signs that tell us rooms 500–519 are to the left
  - And rooms 520–590 are to the right
- With one quick check:
  - We can make the appropriate turn
    - And ignore the rooms in the other direction

The recursive algorithm to find a value

```
FindValue(TreeNode: current, Type: target):
  1. IF current == null:
     return null
  2. IF current.value == target:
     return current
  3. IF target < current.value AND current.left != null:
     return FindValue(current.left, target)
  4. IF target > current.value AND current.right != null:
     return FindValue(current.right, target)
  5. return null
```
Suppose we used this strategy to search for 63
Suppose we used this strategy to search for 63
Suppose we used this strategy to search for 63

The iterative approach to binary search

FindValueItr(TreeNode: root, Type: target):

1. TreeNode: current = root
2. WHILE current != null AND current.value != target:
   3. IF target < current.value:
      current = current.left
   ELSE:
      current = current.right
3. return current
Simplifying the logic for using binary search trees

- We can wrap the entire tree in a **thin** data structure that contains the root node:

  ```
  BinarySearchTree {
    TreeNode: root
  }
  ```

**Growing a Binary Search Tree:**

**Insertion**

It will never rain roses; when we want to have more roses we must plant more trees.

George Eliot
Adding a node to a BST [1/2]

- We use the same basic algorithm to add values to a binary search tree as we do to search it.
- Start at the root node, progress down the tree *as if searching for the new value*, and:
  - Terminate once we hit a dead end:
    - Either a leaf node or an internal node with a single child.

Adding a node to a BST [2/2]

- The primary *difference* between our search and insertion algorithms comes *after we hit the dead end*.
- The insertion algorithm creates a new node as a child of the current node:
  - A left-hand child if the new value is less than that of the current node.
  - A right-hand child if the new value is greater than that of the current node.
For example if we want to add 77

Cost of inserting a node?

- Proportional to the depth of the branch along which we insert the new node
Removing Nodes is more complicated than adding them

- There are three cases of node removal to consider:
  1. Removing a leaf node (with no children)
  2. Removing an internal node with a single child
  3. Removing an internal node with two children
Remove a leaf node: Delete that node & update its parent’s child pointer to reflect that it doesn’t exist

Removal of a leaf node

- Might make the parent node into a leaf
- The value of parent pointers
  - Allows us to follow the parent pointer back to that node’s parent, and set the corresponding child pointer to null
  - Storing this single piece of additional data makes deletion much simpler
Removal: If target node has a single child, **promote that child** to be child of deleted node’s parent.
What about deleting a node with two children [1/2]

- The complexity ramps up substantially
- No longer sufficient to just
  - Delete the node (leaf)
  - Shift a single child up (node with single child)

What about deleting a node with two children [2/2]

- Efficiently find that node’s successor
- While this might seem like a daunting task, it isn’t
- We can always find the successor in the node’s right-hand subtree
  - Specifically, the successor will be the minimum (or leftmost) node in the right-hand subtree
Deleting node 81

Node 81 to be deleted

Deleting node 81
The perils of unbalanced trees

- Binary search trees are still efficient as long as the trees are mostly, if not perfectly, balanced
- But if the tree becomes highly unbalanced?
  - Its depth could grow linearly with the number of elements
- In fact, in the extreme case, our splendid BST becomes nothing more than a sorted linked list
  - All the nodes have a single child in the same direction
Example of a BST devolving into a sorted linked list

Highly unbalanced trees can easily occur in many real-world applications

- Imagine storing coffee log in a BST indexed by timestamp
- Every time we drink a cup of coffee, we insert the relevant information into our tree
- Things go bad quickly
  - Due to the **monotonically increasing timestamps**, we insert every entry in sorted order
  - We end up creating a linked list using only the **right-hand child pointers**
Operations on an unbalanced tree can be extremely inefficient

- Consider a tree with N nodes
- If our tree is balanced?
  - Our operations take time logarithmic in N or $O(\log N)$
- In the opposite case, where our tree is a list?
  - Our operations can take time linearly proportional to N or $O(N)$
- We can use red-black trees, 2-3 trees, and B-trees, to preserve balance
  - While undergoing dynamic insertions and deletions
  - The tradeoff? Increased complexity in the tree operations

**Pivoting to Balance BSTs**
Balancing BSTs

- Perform a rotation step after nodes are added or removed
- If the insert operation leaves a branch unbalanced?
  - i.e., two consecutive nodes in the branch have only one child
  - Rotate nodes around the middle one

Balancing BSTs: Rotation

Pivot node is 3. After rotation, the pivot becomes 5.
B-Trees

- B-Trees combine ideas such as:
  - Increase node fanout
  - Reduce tree height
  - Reduce the number of node pointers
  - Reduce frequency of balancing operations
The contents of this slide-set are based on the following references