# CS 250: Foundation of Computer Systems [Binary Representations \& Operations] 

## A number in context

Look closely ...
a 1 or a 0 in the right place
The leftmost bit to be precise
Is also a sign
Break a sequence of bits here or there
And it gives you powers
To size up who-ville or the universe's atoms
The only clarity that matters? Context
The interpretation that follows? Unambiguous

SHRIDEEP PALLICKARA<br>Computer Science<br>Colorado State University

1

## Frequently asked questions from the previous class

## survey

$\square \mathrm{N}!=1 \times 2 \times 3 \times \ldots \times \mathrm{N}$
$\square$ Did long variables exist on 32-bit systems?
$\square$ Twos complement: Conversion to decimal?

## Topics covered in this lecture

Signed numbers$\square$ Two's complement
$\square$ One's complementFloating point numbersHexadecimal numbers

3

## Announcements

Recitations are moving to CSB-3 15 [Unix Lab]Quiz etiquette
## Summarizing two's complement for an $n$-bit binary system

$\square$ The system codes $2^{n}$ signed numbers, ranging from $-2^{n-1}$ to ( $2^{n-1}-1$ )
$\square$ The code of any nonnegative number begins with a 0
$\square$ The code of any negative number begins with a 1Representation of $-x \rightarrow 2^{n}-x$
$\square$ To obtain the binary code of $-x$ from the binary code of $x$ ?
Flip all the bits of $\boldsymbol{x}$ and add 1 to the result

5

## Let's consider an 8-bit binary system and two's complement

The system codes $2^{n}$ signed numbers, ranging from $-2^{n-1}$ to ( $2^{n-1}-1$ ) i.e., from $-2^{8-1}$ to $\left(2^{8-1}-1\right)$ or $-2^{7}$ to $\left(2^{7}-1\right)$

Let's look at the number -42
Method-1: $-42 \rightarrow 2^{8}-42=214=0 b 11010110$
Method-2: Flip all bits of $x$ and add 1 to the result

- 42 =0b00101010
- Flips the bits: Obl 1010101
- Add 1: Obl 1010110


## Converting 2 s complement to decimal (or denary)

## Obl1010110

$\square$ The weight for the leftmost digit is negative
$\square 1 \times-2^{7}+1 \times 2^{6}+0 \times 2^{5}+1 \times 2^{4}+0 \times 2^{3}+1 \times 2^{2}+1 \times 2^{1}+0 \times 2^{0}$
$\square-128+64+0+16+0+4+2+0$
ロ-128+86
$\square-42_{10}$

## ONE'S COMPLEMENT

Another useful concept to know ... one's complement
$\square$ A not so successful attempt to represent signed numbers
Get to negative numbers by taking positive numbers and flipping all the bits (i.e., 1 becomes 0 and 0 becomes 1)

9

| One's complement | Sign | $2^{2}$ | 21 | $2^{0}$ | Decimal |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 1 | 1 | +7 |
|  | 0 | 1 | 1 | 0 | +6 |
| $\square$ Flipping each bit of 0111 (+7) yields 1000 ( <br> $\square$ Has two different representations for zero | 0 | 1 | 0 | 1 | +5 |
|  | 0 | 1 | 0 | 0 | +4 |
|  | 0 | 0 | 1 | 1 | +3 |
|  | 0 | 0 | 1 | 0 | +2 |
|  | 0 | 0 | 0 | 1 | +1 |
|  | 0 | 0 | 0 | 0 | +0 |
|  | 1 | 1 | 1 | 0 | -1 |
|  | 1 | 1 | 0 | 1 | -2 |
|  | 1 | 1 | 0 | 0 | -3 |
|  | 1 | 0 | 1 | 1 | -4 |
|  | 1 | 0 | 1 | 0 | -5 |
|  | 1 | 0 | 0 | 1 | -6 |
|  | 1 | 0 | 0 | 0 | -7 |
| coldrado state university | Professor: SHRIDEEP PALIICKARA Computer Science Department |  | Binary Representations |  | L4.10 |

10

CS250: Foundations of Computer Systems Dept. Of Computer Science, Colorado State University

## One's complement: Addition is a little more complicated



End-around carry


## General-purpose computers are built to solve general-purpose problems

Which involve a wide range of numbers
$\square$ You can get an idea of this range by skimming a physics textbook
$\square$ There are tiny numbers such as Planck's constant ( $6.63 \times 10^{-34}$ jouleseconds) and
$\square$ Huge numbers such as Avogadro's constant ( $6.02 \times 10^{23}$ molecules/mole)
$\square$ This is a range of $10^{57}$, which comes out to about $2^{191}$

- That's almost 200 bits!

Bits just aren't cheap enough to use a few hundred of them to represent every number, so we need a different approach

## Look to what we have done in the past

Scientific notation represents a large range of numbers by (how else?) creating a new context for interpretation
$\square$ It uses a number with a single digit to the left of the decimal point, called the mantissa, multiplied by 10 raised to some power, called the exponent


## Floating points

$\square$ Computers use the same system, except that the mantissa and exponent are binary numbers and 2 is used instead of 10

This is called the floating point representation

## More on the scientific notation: Let's come up with a

 representation$$
\pm \square . \square \mathrm{e} \pm \square
$$

A value like 9,876,543,210 would be approximated with $9.88 \times 10^{9}$ (or $9.88 \mathrm{e}+9$ in programming language notation)

## Scientific notation complicates arithmetic somewhat

$\square$ When adding and subtracting two numbers in scientific notation, you must adjust the two values so that their exponents are the same
$\square$ For example, when adding 1.23 e 1 and 4.56 e 0 , you could convert 4.56 e 0 to 0.456 e 1 and then add them
$\square$ The result ( $1.23 \mathrm{e} 1+0.456 \mathrm{e} 1=1.686 \mathrm{e} 1)$, does not fit into the three significant digits of our current format (in the previous slide)
$\square$ So, we must either round or truncate the result to three significant digits - Rounding generally produces the more accurate result, so let's round to obtain 1.69el

## The implications of such adjustments?

The lack of precision (the number of digits or bits maintained in a computation) affects the accuracy (the correctness of the computation)

## Floating point representation: Binary version of the scientific notation to represent a wide range of numbers

$\square$ The naming convention is confusing because the binary (or decimal) point is always in the same place
$\square$ Between the ones and halves (tenths in decimal)

- The "float" is just another way of saying "scientific notation," which allows us to write $1.2 \times 10^{-3}$ instead of 0.0012

By separating the significant digits from the exponents
$\square$ The floating-point system allows us to represent very small or very large numbers without having to store all those zeros

## IEEE 754 Floating standard

When Intel planned to introduce a floating-point unit (FPU) for its original 8086 microprocessor?
$\square$ Intel knew their electrical engineers and solid-state physicists didn't have the numerical analysis background to design a good floating-point representation
Went out and hired the best numerical analyst they could find to design a floatingpoint format for its 8087 FPU
$\square$ That person then hired two other experts in the field, and the three of them (Kahan, Coonen, and Stone) designed the KCS Floating-Point Standard
$\square$ They did such a good job that the IEEE organization used this format as the basis for the IEEE Std 754 floating-point format

The floating-point system is the standard way to represent real numbers in computingThere are two signs:
$\square$ One for the mantissa and one for the exponent (hidden)
$\square$ There are also a lot of tricks to make sure that things like rounding work as well as possible and to minimize the number of wasted bit combinations

## The mantissa and the exponent

$\square$ The mantissa is a base value that usually falls within a limited range (for example, between 0 and 1)
$\square$ The exponent is a multiplier that, when applied to the mantissa, produces values outside this range
$\square$ The big advantage of the mantissa/exponent configuration $\square$ Floating-point format can represent values across a wide range

## IEEE Floating point number formats



## A couple of the tricks that are used in the standard

$\square$ Normalization, which adjusts the mantissa so that there are no leading (that is, on the left) zeros

A second trick, from Digital Equipment Corporation (DEC), doubles the accuracy
$\square$ By throwing away the leftmost bit of the mantissa since we know that it will always be 1, which makes room for one more bit

Exponent values of all $0 s$ and all 1 s would have special meaning

## Two types of floating point numbers: Single and double-precision

Single-precision numbers use 32 bits and can represent numbers approximately in the range $\pm 10^{ \pm 38}$ with about 7 digits of accuracy
$\square$ Although there is an infinite number of values between 1 and 2, we can represent only 8 million ( $2^{23}$ ) of them because we use a 23 -bit mantissa - Therefore, have only 23 bits of precision

Double-precision numbers use 64 bits and can represent a wider range of numbers, approximately $\pm 10 \pm 308$, with about 15 digits of accuracy
Number of atoms in the known universe is between $10^{78} \sim 10^{82}$

## IEEE 754 also uses some special bit patterns

$\square$ To represent things like division by zero, which evaluates to positive or negative infinity

It also specifies a special value called NaN , which stands for "not a number"

- If you find yourself in the NaNny state, it probably means that you did some illegal arithmetic operation


## Let's look at a number with a decimal point

$\square 0.15625$We will compute its binary representation


## Unlike the division (by 2) that we did for number left of the decimal point ....

During conversions, for numbers to the right of the decimal, we will multiply 2
$\square 0.15625_{10}$


| $0.15625 \times 2=0.3125$ | 0 |
| :--- | :--- |
| $0.3125 \times 2=0.625$ | 0 |
| $0.625 \times 2=1.250$ | 1 |
| $0.250 \times 2=0.500$ | 0 |
| $0.500 \times 2=1.000$ | 1 |

## $0.15625_{10}=0.00101_{2}$

$\square 0.00101_{2}=1.01 \times 2^{-3}$
$\square$ Fraction is .01 and the exponent is -3The number is positive
$\square$ In the IEEE-754 standard, the exponent is written as a biased exponent In single-precision, you need to add 127

- -3 would be written as $-3+127=124_{10}=01111100_{2}$

In double-precision you need to add 1023

## Representation of $0.15625_{10}=0.00101_{2}$ in IEEE- <br> 754 single precision

$0.00101_{2}=1.01 \times 2^{-3}$The number is positiveExponent is written as a biased exponent: $01111100_{2}$We only write the fractional part (the 1 in 1.01 is implied) i.e., 01 and then pad zeros all the way to the right

## NaN and Infinity

The biased-exponent field is filled with all 1 bits to indicate either- Infinity mantissa field $=0$
$\square \mathrm{NaN}$ mantissa field $\neq 0$


32

## Hexadecimal

$\square$ Hexadecimal is base 16!
$\square$ Given what we've already seen so far, you probably know what that meansHexadecimal, or just hex for short, is a place-value system where
Each place represents a power of 16
$\square$ and each place can be one of 16 symbols

## Hexadecimal number representation

As in all place-value systems, the rightmost place will still be the ones place
$\square$ The next place to the left will be the sixteens place, then the 256 s (16 $\times 16)$ place, then the 4,096 s $(16 \times 16 \times 16)$ place, and so on

Simple enough!

## But what about the other requirement that each place can be one of 16 symbols?

We usually have ten symbols to use to represent numbers, 0 through 9
$\square$ We need to add six more symbols to represent the other values

But what about the other requirement that each place can be one of 16 symbols?
$\square$ We could pick some random symbols like \& @ \#, but these symbols have no obvious order
$\square$ Instead, the standard is to use $A, B, C, D, E$, and $F$
Either uppercase or lowercase is fine!
$\square$ In this scheme, A represents ten, B represents eleven, and so on, up to F, which represents fifteen


37

## Consider the number $0 \times 1$ A5 in hexadecimal

$\square$ What's the value of this number in decimal?
$\square$ The rightmost place is worth 5The next place has a weight of 16 , and there's an A there, which is 10 in decimal, so the middle place is worth $16 \times 10=160$The leftmost place has a weight of 256 , and there's a 1 in that place, so that place is worth 256The total value then is $256+160+5=421$ in decimal


## Conversion from decimal to hex

$\square$ The hex number is constructed from right to left
$\square$ Step 1: Divide the decimal number by 16 and note down the remainder
$\square$ Step 2: Divide the obtained quotient by 16, and note remainder again
$\square$ Step 3: Repeat the above steps until you get 0 as the quotient Stopping criteria

## Example: Decimal number 421

$\square 421 \div 16=26 \quad$ (Remainder 5)$26 \div 16=1 \quad$ (Remainder 10 or $\mathbf{A})$$1 \div 16=0 \quad$ (Remainder 1)

Hex representation: 0x1A5
口 $=1 * 256+10 * 16+1 * 5$
$\square=256+160+5$
$\square=421$ in decimal

## Some conversions across number systems

|  | Example 1 | Example 2 |
| :--- | :--- | :--- |
| Binary | 1111000000001111 | 1000100010000001 |
| Hexadecimal | FOOF | 8881 |
| Decimal | 61,455 | 34,945 |



43

## Why binary logic matters

$\square$ Every digital device - be it a PC, a cell phone, or a network router - is based on chips designed to store and process binary information
$\square$ Although these chips come in different shapes and forms, they are all made of the same building blocks: elementary logic gates
$\square$ The gates can be physically realized using many different hardware technologies

But their logical behavior, or abstraction, is consistent across all implementations

## We've looked at using binary to represent data, but computers do more than simply store data

$\square$ Binary allows us to work with data as well
$\square$ Computers give us the capability to process data using hardware that we can program to execute a sequence of simple instructions

- Instructions like "add two numbers together" or "check if two values are equal"

Computer processors that implement these instructions are fundamentally based on binary logic
$\square$ A system for describing logical statements where variables can only be one of two values - true or false

## Why is binary a natural fit for logic?

$\square$ Typically, when someone speaks of logic, they mean reasoning, or thinking through what is known in order to arrive at a valid conclusion

When presented with a set of facts, logic allows us to determine whether another related statement is also factual
$\square$ Logic is all about truth - what is true, and what is false
Likewise, a bit can only be one of two values, 1 or 0
$\square$ Therefore, a single bit can be used to represent a logical state of true (1) or false (0)

## Let's consider the logical statements for a rectangle

$\square$ If the shape does not have four sides and does not have four right angles
$\square \mathrm{lt}$ is not a rectangle
$\square$ If the shape does not have four sides and does have four right angles
$\square$ It is not a rectangle
$\square$ If the shape does have four sides and does not have four right angles
$\square$ It is not a rectangle
$\square$ If the shape does have four sides and does have four right angles
$\square$ It is a rectangle!

## Let's put those statements in table

| Four sides | Four right angles | Is a rectangle |
| :---: | :--- | :--- |
| False | False | False |
| False | True | False |
| True | False | False |
| True | True | True |

## What's so special about this table?

This type of table is known as a truth tableA truth table shows all the possible combinations ofConditions (inputs) and
$\square$ Their logical conclusions (outputs)Our previous table was written specifically for our statement about a rectangle, but ...
$\square$ The same table applies to any logical statement joined with AND

## Let's represent false as 0 and true as a 1

| $x$ | $y$ | $x$ And $y$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

## Let's look at another example: OR

Say you work at a shop that gives a discount to only two types of customers:
$\square$ Children

- People wearing sunglasses
$\square$ No one else is eligible for a discount
$\square$ If you wanted to state the store's policy as a logical expression, you could say the following:
$\square$ GIVEN the customer is a child
$\square$ OR GIVEN the customer is wearing sunglasses
$\square$ I CONCLUDE that the customer is eligible for a discount


## The OR Truth Table

| $x$ | $y$ | $x$ Or $y$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |



53

## Logic Operations

$\square$ One use of bits is to represent the answers to yes/no questions such as "Is it cold?" or "Do you like my hat?"
$\square$ We use the terms true for yes and false for no
$\square$ Questions like "Where's the party?" don't have a yes/no answer and can't be represented by a single bit
$\square$ We often combine several yes/no clauses into a single sentence $\square$ We might say, "Wear a coat if it is cold or if it is raining" or "Go skiing if it is snowing and it's not a school day"

## Logic Operations

Another way of saying those things might be
口 "Wear coat is true if cold is true or raining is true" and "Skiing is true if snowing is true and school day is not true"
$\square$ These are logic operations that each produce a new bit based on the contents of other bits

## Boolean Functions

$\square$ A boolean function is a function that operates on binary inputs and returns binary outputs

Since computer hardware is based on representing and manipulating binary values ...
$\square$ Boolean functions play a central role in the specification, analysis, and optimization of hardware architectures

## Every Boolean function can be defined using two alternative representations

$\square$ First, we can define the function using a truth table
For each one of the $2^{n}$ possible tuples of input variable values, the table lists the value of $f\left(v_{1}, v_{2}, \ldots, v_{n}\right)$
$\square$ Can be thought of as a data-driven definition
In addition to this data-driven definition, we can also define Boolean functions using Boolean expressions
$\square$ For example: (x Or y) And Not (z)

## Boolean Algebra

Algebra is a set of rules for operating on numbers
Boolean algebra manipulates two-state binary values that are typically labeled true/false, 1/0, yes/no, on/off, and so forth

We will use 1 and 0

## Boolean Algebra

Boolean algebra, invented in the 1800s by English mathematician George Boole, is a set of rules that we use to operate on bits

As with regular algebra: the associative, commutative, and distributive rules also apply
$\square x{ }^{*} y=y$ * $x \quad$ Commutative
$\square(x * y) * z=x *(y * z) \quad$ Associative
$\square x *(y+z)=x^{*} y+x^{*} z \quad$ Distributive

## The contents of this slide-set are based on the following references

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