

Transistor: Building Block of Computers

- Microprocessors contain lots of transistors
- Intel 8086 (1978): 29 thousand
- Intel 80186 (1982): 55 thousand
- Intel 80386 (1985): 275 thousand
- Intel 80486 (1989): 1.1 million
- Intel Pentium (1993): 3.1 million
- Intel Pentium II (1998): 7.5 million
- Intel Pentium III (2001): 45 million
- Intel Pentium 4 (2006): 184 million
- Intel Core 2 Duo (2006): 291 million
- Intel Quad Core i7 (2011): 1.1 billion
- Intel 8-core Xeon (2012): 2.3 billion


## Transistor: Building Block of Computers

- Logically, each transistor acts as a switch
- Combined to implement logic functions (gates)
- AND, OR, NOT
- Combined to build higher-level structures
- Multiplexer, decoder, register, memory ...
- Adder, multiplier ...
- Combined to build simple processor
- LC-3




## 

## PropagationDelay

- Each gate has a propagation delay, typically fraction of a nanosecond ( $10^{-9} \mathrm{sec}$ ).
- Delays accumulate depending on the chain of gates the signals have to go through.
- Clock frequency of a processor is determined by the delay of the longest combinational path between storage elements, i.e. cycle time.



## Logical Operation: OR and NOR






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## BooleanAlgebraLaws

- Commutative
- $A+B=B+A$
- $A \cdot B=B \cdot A$
- Associative
- $A+(B+C)=(A+B)+C=A+B+C$
- $A \cdot(B \cdot C)=(A \cdot B) \cdot C=A B C$
- Distributive
- $A \cdot(B+C)=A \cdot B+A \cdot C$
- $A+(B \cdot C)=(A+B) \cdot(A+C)$



## Some Useful Identities for simplification

## - $A \bar{B}+A B=A$

Proof: $\overline{A B}+A B=\overline{A(\bar{B}}+B) / /$ Distributive Law

| $=A(T)$ | // Negation Law |
| :--- | :--- |
| $=A$ | I/ Identity Law |

- $A+A B=A$

Proof: $A+A B=A(1+B)$ // Distributive Law = A(1) // Domination Law = A I/ Identity Law

## DeMorgan's Law

- Converting AND to OR (with some help from NOT) - Consider the following gate:


| $A$ | $B$ | $\bar{A}$ | $\bar{B}$ | $\bar{A} \cdot \bar{B}$ | $\overline{\mathrm{~A}} \cdot \overline{\mathrm{~B}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 1 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 | 1 |

To convert AND to OR (or vice versa), invert inputs and output.
Same as A OR B!

## copmg end

## More than 2 Inputs?

- AND/OR can take any number of inputs.
- AND = 1 if all inputs are 1 .
- $O R=1$ if any input is 1 .
- Similar for NAND/NOR.
- Can implement with multiple two-input gates, or with single CMOS circuit.


Summary
- MOS transistors are used as switches to implement logic functions.
- n-type: connect to GND, turn on (1) to pull down to 0
- p-type: connect to +2.9 V , turn on ( 0 ) to pull up to 1
- Basic gates: NOT, NOR, NAND
- Logic functions are usually expressed with AND, OR, and NOT
- DeMorgan's Law
- Convert AND to OR (and vice versa) by inverting inputs and output


## BuildingFunctions from Logic Gates

- Combinational Logic Circuit
- output depends only on the current inputs
- stateless
- Sequential Logic Circuit
- output depends on the sequence of inputs (past and present)
- stores information (state) from past inputs
- We'll first look at some useful combinational circuits, then show how to use sequential circuits to store information.

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[^0]:    CS270-Sping Semester 2016

