

## Chapter 3 Digital Logic Structures

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## Computing Layers

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## Combinational vs. Sequential

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- **Combinational Circuit**
  - does not store information, always gives the same output for a given set of inputs
    - *example:* adder always generates sum and carry, regardless of previous inputs
- **Sequential Circuit**
  - stores information, output depends on stored info (state) plus input
  - so a given input might produce different outputs, depending on the stored information
  - useful for building “memory” elements and “state machines”
    - *example:* ticket counter

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## R-S Latch: Simple Storage Element

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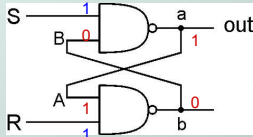
- R is used to “reset” or “clear” the element – set it to zero.
- S is used to “set” the element – set it to one.

- If both R and S are one, output could be either zero or one.
  - “quiescent” state -- holds its previous value
  - if a is 1, b is 0, and vice versa

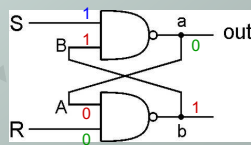
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### Clearing the R-S latch

- Suppose we start with output = 1, then change R to zero.



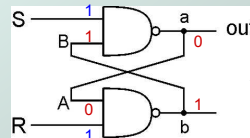
Output changes to zero.



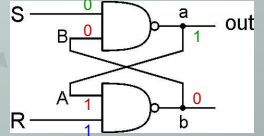
Then set R=1 to "store" value in quiescent state.

### Setting the R-S Latch

- Suppose we start with output = 0, then change S to zero.



Output changes to one.



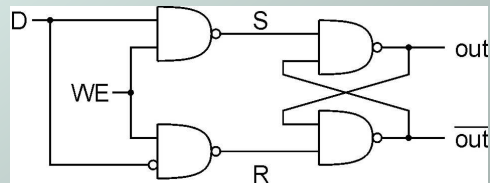
Then set S=1 to "store" value in quiescent state.

### R-S Latch Summary

- R = S = 1**
  - hold current value in latch
- S = 0, R=1**
  - set value to 1
- R = 0, S = 1**
  - set value to 0
- R = S = 0**
  - both outputs equal one
  - final state determined by electrical properties of gates
  - Don't do it!**

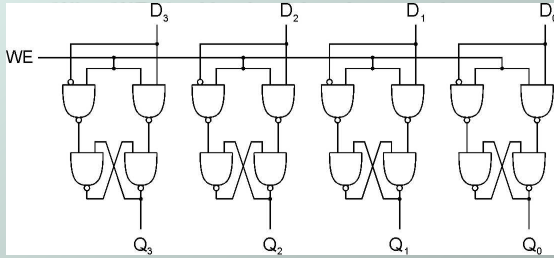
### Gated D-Latch

- Two inputs: D (data) and WE (write enable)
  - when **WE = 1**, latch is set to **value of D**
    - $S = \text{NOT}(D), R = D$
  - when **WE = 0**, latch holds **previous value**
    - $S = R = 1$



### Register

- A register stores a multi-bit value.
  - We use a collection of D-latches, all controlled by a common WE.



### Representing Multi-bit Values

- Number bits from right (0) to left (n-1)
  - just a convention -- could be left to right, but must be *consistent*

- Use brackets to denote range:
  - $D[l:r]$  denotes bit l to bit r, from *left to right*

$$A = \overset{15}{0}1010011010101\overset{0}{1}$$

$$A[14:9] = 101001$$

$$A[2:0] = 101$$

- May also see  $A\langle 14:9 \rangle$ , especially in hardware block diagrams.

### Memory

- Now that we know how to store bits, we can build a memory – a logical  $k \times m$  array of stored bits.

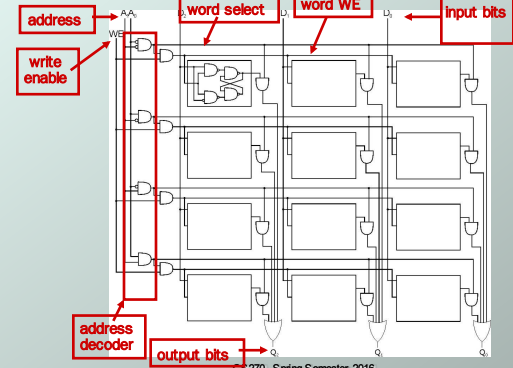
**Address Space:**  
number of locations  
(usually a power of 2)

$$k = 2^n \text{ locations}$$

**Addressability:**  
number of bits per location  
(e.g., byte-addressable)

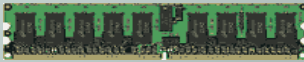


### $2^2 \times 3$ Memory



## More Memory Details

- ◆ Not the way actual memory is implemented!
  - fewer transistors, denser, relies on electrical properties
- ◆ But the logical structure is very similar.
  - address decoder, word select line, word write enable
- ◆ Random Access Memory: 2 different types
  - **Static RAM** (SRAM)
    - fast, used for caches, maintains data when powered
  - **Dynamic RAM** (DRAM)
    - slower but denser, storage decays, must be refreshed
- ◆ Non-Volatile Memory: **ROM, PROM, Flash**



## Memory Bandwidth

- ◆ Bandwidth is the rate at which memory can be read or written by the processor.
- ◆ Approximately equal to the memory bus size times the speed at which the memory is clocked.
- ◆ Examples of bandwidth (from Wikipedia):
  - Phone line, Modem, up to 5.6KB/s
  - Digital subscriber line, ADSL, up to 128KB/s
  - Wireless networking, 802.11g, up to 17.5MB/s
  - Peripheral connection, USB 2.0, 60MB/s
  - Digital video, HDMI, up to 1.275GB/s
  - Computer bus, PCI Express, up to 25.6GB/s
  - Memory chips, SDRAM, up to 52GB/s

## Looking Ahead: C Arrays

- ◆ Similar to Java arrays

```
// integer array
int iArray[3] = {1,2,3};
printf("iArray[2]: %d", iArray[2]);

// float array
float fArray[2] = {0.1f,0.2f};
printf("fArray[1]: %f", fArray[1]);

// character array
char cArray[4] = {'a','b','c','d'};
printf("cArray[3]: %c", cArray[3]);
```

## Looking Ahead: C Strings

- ◆ Array of chars with null (not NULL) termination

```
// string: static allocation
char *string1 = "Hello World\n";
printf("string1: %s", string1);

// string: dynamic allocation
char *string2 = malloc(13);
strcpy(string2, "Hello World\n");
```

Note that the programmer is responsible for making sure string has enough memory!

## Looking Ahead: C Arrays and C Pointers

- ◆ Array name is a pointer to array

```
int iArray[2] = {1234, 5678};

printf("iArray[0]: %d", iArray[0]);
printf("iArray[1]: %d", iArray[1]);
printf("&iArray[0]: %x", &iArray[0]);
printf("&iArray[1]: %x", &iArray[1]);
printf("iArray: %x", iArray);
iArray[2] = 0; // out of bounds!
```

## Looking Ahead: C Functions

- ◆ Can pass by value or reference

```
// by value (copies value)
float f1(int i, float f);
// by reference (copies pointer)
float f2(float *f);
```

- ◆ Function cannot change values passed by value

```
f1: i = 10; // changes the copy
```

- ◆ Function can change values passed by reference

```
f2: *f = 1.2; // changes actual value
```