Transformations: Narrow and Wide
Though their numbers are few
Don’t let them beguile you
Innocuous though they may seem
The wrong invocation is all it takes
To amplify inefficiencies And protract computations

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Topics covered in this lecture
- Pair RDDs
- Data Frames
- Dependencies and Transformations
- Partitioners
Transformations on Pair RDDs

Pair RDDs

- RDDs that contain **key/value pairs**
- Expose partitions that allow you to act on each key in parallel or regroup data across the network
Transformations on Pair RDDs

- Pair RDD = \{(1,2), (3,4), (3,6) \}

- `reduceByKey(func)`
  - Combine values with the same key
  - Invocation: `rdd.reduceByKey((x, y) => x + y)`
  - Result: \{(1, 2), (3, 10) \}

- `groupByKey(func)`
  - Group values with the same key
  - Invocation: `rdd.groupByKey()`
  - Result: \{(1, [2]), (3, [4, 6]) \}
Transformations on Pair RDDs [3/5]

- Pair RDD = \{(1,2), (3,4), (3,6) \}

- `mapValues(func)`
  - Apply function to each value of a pair RDD *without* changing the key
  - Invocation: `rdd.mapValues(x => x+1)`
  - Result: \{(1, 3), (3, 5), (3, 7) \}

Transformations on Pair RDDs [4/5]

- Pair RDD = \{(1,2), (3,4), (3,6) \}

- `values()`
  - Return an RDD of just the values
  - Invocation: `rdd.values()`
  - Result: \{ 2, 4, 6 \}
Transformations on Pair RDDs

- Pair RDD = \{(1,2), (3,4), (3,6) \}

- `sortByKey()`
  - Return an RDD sorted by the key
  - Invocation: `rdd.sortByKey()`
  - Result: \{(1,2), (3,4), (3,6) \}
Transformations on two Pair RDDs

- **subtractByKey()**
  - Remove elements with a key present in the other RDD
  - Invocation: `rdd.subtractByKey(other)`
  - Result: `{(1,2)}`

- **join()**
  - Perform an inner join between two RDDs. Only keys that are present in both pair RDDs are output
  - Invocation: `rdd.join(other)`
  - Result: `{(3, (4,9)), (3, (6,9))}`
Transformations on two Pair RDDs

- \( rdd = \{(1,2), (3,4), (3,6)\} \quad \text{other} = \{(3,9)\} \)

- **leftOuterJoin()**
  - Perform a join between two RDDs where the **key must be present in the first RDD**
  - Value associated with each key is a tuple of the value from the source and an Option for the value from the other pair RDD
    - In python if a value is not present, None is used.
  - Invocation: \( rdd.leftOuterJoin(\text{other}) \)
  - Result: \{ (1, (2,None)) , (3, (4, 9)) ,  (3, (6, 9)) \}

- **rightOuterJoin()**
  - Perform a join between two RDDs where the key must be present in the other RDD;
  - Tuple has an option for the source rather than other RDD
  - Invocation: \( rdd.rightOuterJoin(\text{other}) \)
  - Result: \{ (3, (4,9) ) , (3, (6,9)) \}
Transformations on two Pair RDDs

- \( \text{rdd} = \{(1,2), (3,4), (3,6)\} \quad \text{other} = \{(3,9)\} \)

- \( \text{cogroup()} \)
  - Group data from both RDDs using the same key
  - Invocation: \( \text{rdd}.\text{cogroup}(\text{other}) \)
  - Result: \( \{(1, ([2], [])), (3, ([4, 6], [9]))\} \)

Example of chaining operations:
Calculation of per-key average

\[
\begin{array}{|c|c|}
\hline
\text{key} & \text{value} \\
\hline
\text{panda} & 0 \\
\text{pink} & 3 \\
\text{pirate} & 3 \\
\text{panda} & 1 \\
\text{pink} & 4 \\
\hline
\end{array}
\]

\[
\begin{array}{|c|c|}
\hline
\text{key} & \text{value} \\
\hline
\text{panda} & (1, 2) \\
\text{pink} & (7, 2) \\
\text{pirate} & (3, 1) \\
\hline
\end{array}
\]

\[
\begin{array}{|c|c|}
\hline
\text{key} & \text{value} \\
\hline
\text{panda} & 0 \quad (0, 1) \\
\text{pink} & 3 \quad (3, 1) \\
\text{pirate} & 3 \quad (3, 1) \\
\text{panda} & 1 \quad (1, 1) \\
\text{pink} & 4 \quad (4, 1) \\
\hline
\end{array}
\]

\[
\text{rdd}.\text{mapValues}(x => (x, 1)).\text{reduceByKey}( (x, y) => (x._1 + y._1, x._2 + y._2))
\]
A word count example

- We are using `flatMap()` to produce a pair RDD of words and the number 1

```python
rdd = sc.textfile("s3://...")
words = rdd.flatMap(lambda x: x.split(" "))
result = words.map(lambda x: (x, 1)).reduceByKey(lambda x, y: x+y)
```
Spark DataFrame

- DataFrames consist of
  - A series of records (like rows in a table) that are of type Row
  - A number of columns (like columns in a spreadsheet)

- Rows
  - You can create rows by manually instantiating a Row object with the values that belong in each column

- Columns
  - You can select, manipulate, and remove columns from DataFrames and these operations are represented as expressions

Schemas

- A schema defines the column names and types of a DataFrame
- You can let a data source define the schema (called schema-on-read) or define it explicitly
- Note that only DataFrames have schemas
  - Rows themselves do not have schemas.
  - If you create a Row manually?
    - You must specify the values in the same order as the schema of the DataFrame to which they might be appended
We can create DataFrames from raw data sources

- Spark has six "core" data sources
  - CSV
  - JSON
  - Parquet
  - ORC: Apache Optimized Row Columnar (ORC) file format
  - JDBC/ODBC connections
  - Plain-text files

- Hundreds of external data sources written by the community
  - E.g.: Cassandra, HBase, MongoDB, AWS, Redshift, XML etc.

The foundation for reading data in Spark is the DataFrameReader

- We access this through the SparkSession via the read attribute: spark.read
- After we have a DataFrame reader, we specify several values:
  - The format: Input data source format
  - The schema
  - The read mode {Permissive, DropMalformed, Failfast}
  - A series of options
- The format, options, and schema each return a DataFrameReader that can undergo further transformations and are all optional
However, at a minimum, the DataFrameReader must have a **path** from which to read

\[
\text{spark.read.format("csv")}
\]
\[
\quad .\text{option("mode", "FAILFAST")}
\]
\[
\quad .\text{option("inferSchema", "true")}
\]
\[
\quad .\text{option("path", "path/to/file(s)")}
\]
\[
\quad .\text{schema(someSchema)}
\]
\[
\quad .\text{load()}
\]

Writing data is quite similar to that of reading data

- Instead of the DataFrameReader, we have the DataFrameWriter
- We access the DataFrameWriter on a per-DataFrame basis via the write attribute:

\[
dataFrame.write
\]
Writing Data

- After we have a DataFrameWriter, we specify three values:
  - The format, a series of options, and the save mode
- At a minimum, you must supply a path.
- Options may vary from data source to data source.

```scala
df.write.format("csv")
  .option("mode", "APPEND")
  .option("dateFormat", "yyyy-MM-dd")
  .option("path", "path/to/file(s)")
  .save()
```

You can make any DataFrame into a table or view

- Done via a simple method call: `createOrReplaceTempView`
- This then allows you to query the data using SQL

```scala
val df = spark.read
  .format("json")
  .load("/data/flight-data/json/2015-summary.json")

df.createOrReplaceTempView("dfTable")
```
DataFrame transformations

- Add rows or columns
- Remove rows or columns
- Transform a row into a column (or vice versa)
- Change the order of rows based on the values in columns

Adding Columns

- Use the `withColumn` method on the DataFrame
- For example, let's add a column that just adds the number one as a column:

  ```scala
df.withColumn("numberOne", lit(1))
  ```
Renaming Columns

- Done using the `withColumnRenamed` method.
- Will rename the column with the name of the string in the first argument to the string in the second argument:

```java
df.withColumnRenamed("DEST_COUNTRY_NAME","dest")
```

Removing Columns

- Done using a method called `drop`
  ```java
df.drop("ORIGIN_COUNTRY_NAME")
```
- We can drop multiple columns by passing in multiple columns as arguments
  ```java
dfWithLongColName.drop("ORIGIN_COUNTRY_NAME", "DEST_COUNTRY_NAME")
```
Filtering Rows

- To filter rows, we create an expression that evaluates to true or false.
  - Those rows where the expression evaluates to false are filtered out.

  \[ \text{df.filter( col( "count" ) < 2) } \]

Getting Unique Rows

- A very common use case is to extract the unique or distinct values in a DataFrame.
  - These values can be in one or more columns.
  - Done by using the distinct method on a DataFrame.
    - Allows deduplication of any rows that are in that DataFrame.
  - Again, this is a transformation that will return a new DataFrame with only unique rows:

  \[ \text{df.select("ORIGIN_COUNTRY_NAME","DEST_COUNTRY_NAME") .distinct( )} \]
Random Samples

- You might want to sample some random records from a DataFrame
- Done by using the `sample` method on a DataFrame
  - Specify a fraction of rows to extract from a DataFrame and whether the sample will be with or without replacement

```scala
val seed = 5
val withReplacement = false
val fraction = 0.5
df.sample(withReplacement, fraction, seed)
```

Random Splits

- Random splits are helpful when you need to break up a DataFrame into a random “splits” of the original DataFrame
- Often used with machine learning algorithms to create training, validation, and test sets

```scala
val dataFrames =
    df.randomSplit(Array (0.25, 0.75 ), seed)
```
Column Manipulations

- `withColumn(columnName, func)`
  - Return an DataFrame with the additional column
  - Invocation: `df.withColumn("dogYears", df.age / 7)`

- `dropColumn(columnName)`
  - Return an DataFrame without the column
  - Invocation: `df.dropColumn("age")`

Column Manipulations

- `select(columnNames)`
  - Return an DataFrame with the specified columns
  - Invocation: `df.select("firstName", "age")`

- `describe(columnName)`
  - Compute summary statistics over DataFrame columns
  - Invocation: `df.describe("age")`, `df.describe()`
val df = Seq(
    ("Peterson", "Marcus", 54),
    ("Batey", "Edward", 36),
    ("Bruce", "Karen", 35)
).toDF("lastName", "firstName", "age")

df.withColumn("dogYears", df.age / 7.0)
df.describe("age", "dogYears")

+-------+--------+----------+
| summary|  age   |  dogYears|
|--------+--------+----------+
| count  |  3     |  3       |
| mean   | 41.6667|  5.95238 |
| stddev | 10.69268|  1.52753 |
| min    |  35    |   5      |
| max    |  54    |  7.714286|
+-------+--------+----------+
Dataframe joins

- `join(other, <columnComparison>, <joinType>)`
  - Performs a join between 2 Dataframes
  - Invocation: `df1.join(df2, Seq("id"))`

Join column comparison

- Supports a variety of criteria
  - Sequence of column names (e.g. `Seq("id", "age")`)
  - Elaborate comparison definitions (e.g. `df1("age") >= df2("age")`)
Join Type

- DataFrames may perform multiple styles of join
  - Inner: typical dataset join with key to key match
  - Outer, left-outer, right-outer: result contains all rows, filling in columns with 'null' values where data doesn't exist
  - Left-semi, right-semi: similar to outer join, but result only contains rows in specified source dataset

Example: Spark SQL

```scala
val df = Seq(
  ("Peterson", "Marcus", 54),
  ("Batey", "Edward", 36),
  ("Bruce", "Karen", 35)
).toDF("lastName", "firstName", "age")

df.createOrReplaceTempView("people")
spark.sql("SELECT firstName, age, age / 7.0 as dogYears FROM people where age < 50")
```
Tuning the level of parallelism

- Every RDD has a **fixed number of partitions**
  - Determine the degree of parallelism when executing operations

- During aggregations or grouping operations, you can ask Spark to use a specific number of partitions
  - This will override defaults that Spark uses
Example: Tuning the level of parallelism

```python
data = ["a", 3), ("b", 4), ("a", 1)]

sc.parallelize(data).
   reduceByKey(lambda x, y: x + y) #default

sc.parallelize(data).
   reduceByKey(lambda x, y: x + y, 10) #Custom
```

What if you want to tune parallelism outside of grouping and aggregation operations?

- There is `repartition()`
  - **Shuffles data across the network** to create a new set of partitions
  - Very expensive operation!

- There is the `coalesce()` operation
  - Allows avoiding data movement
    - But only if you are **decreasing** the number of partitions
  - Check `rdd.getNumPartitions()` and make sure you are coalescing to fewer partitions than current
Transformations and Dependencies

- Two categories of **dependencies**
  - Narrow
    - Each partition of the parent RDD is used by **at most one partition of the child** RDD
  - Wide
    - Multiple child RDD partitions may depend on a single parent RDD partition

- The narrow versus wide distinction has significant implications for the way Spark evaluates a transformation and, consequently, for its performance
Narrow Transformations

- Narrow transformations are those in which each partition in the child RDD has simple, finite dependencies on partitions in the parent RDD.
- Can be determined at design time, irrespective of the values of the records in the parent partitions.
- Partitions in narrow transformations can either depend on:
  - One parent (such as in the map operator), or
  - A unique subset of the parent partitions that is known at design time (coalesce).
- Narrow transformations can be executed on an arbitrary subset of the data without any information about the other partitions.

Dependencies between partitions for narrow transformations

Diagram showing dependencies between parent and child partitions.
Wide Transformations

- Transformations with **wide dependencies** cannot be executed on arbitrary rows
- Require the data to be partitioned in a particular way, e.g., according the value of their key
  - In sort, for example, records have to be partitioned so that keys in the same range are on the same partition
- Transformations with wide dependencies include sort, reduceByKey, groupByKey, join, and anything that calls the rePartition function

Dependencies between partitions for wide transformations

Wide dependencies **cannot be known fully before the data is evaluated**

The dependency graph for any operations that cause a shuffle (such as groupByKey, reduceByKey, sort, and sortByKey) follows this pattern
PAIR RDDs: What to Watch For

Despite their utility, key/value operations can lead to a number of performance issues

- Most expensive operations in Spark fit into the key/value pair paradigm
  - Because most wide transformations are key/value transformations,
    - And most require some fine tuning and care to be performant
In particular, operations on key/value pairs can cause …

1. Out-of-memory errors in the driver
2. Out-of-memory errors on the executor nodes
3. Shuffle failures
4. “Straggler tasks” or partitions, which are especially slow to compute

The last three performance issues are all most often caused by shuffles associated with the wide transformations.

Memory errors in the driver, are usually caused by actions

- Several key/value actions (including countByKey, countByValue, lookUp, and collectAsMap) return data to the driver
- In most instances they return unbounded data since the number of keys and the number of values are unknown
- In addition to number of records, the size of each record is an important factor in causing memory errors
Preventing out-of-memory errors with aggregation operations [1/2]

- `combineByKey` and all of the aggregation operators built on top of it (`reduceByKey`, `foldLeft`, `foldRight`, `aggregateByKey`) may lead to memory errors if they cause the accumulator to become too large for one key.

- What about `groupByKey`?
  - It is actually implemented using `combineByKey` where the accumulator is an iterator with all the data.

Preventing out-of-memory errors with aggregation operations [2/2]

- Use functions that implement **map-side combinations**
  - Meaning that records with the same key are combined before they are shuffled.
  - This can greatly reduce the shuffled read.

- The following four functions are implemented to use map-side combinations:
  - `reduceByKey`
  - `treeAggregate`
  - `aggregateByKey`
  - `foldByKey`
Two primary techniques to avoid performance problems associated with shuffles

- Shuffle Less
- Shuffle Better

Shuffle Less

- Preserve partitioning across narrow transformations to avoid reshuffling data
- Use the same partitioner on a sequence of wide transformations. This can be particularly useful:
  - To avoid shuffles during joins and …
  - To reduce the number of shuffles required to compute a sequence of wide transformations
Shuffle Better [1/2]

- Sometimes, computation cannot be completed without a shuffle
- However, not all wide transformations and not all shuffles are equally expensive or prone to failure

Shuffle Better [2/2]

- By using wide transformations such as `reduceByKey` and `aggregateByKey` that can perform map-side reductions and that do not require loading all the records for one key into memory?
  - You can prevent memory errors on the executors and
  - Speed up wide transformations, particularly for aggregation operations
- Lastly, shuffling data in which records are distributed evenly throughout the keys, and which contain a high number of distinct keys?
  - Prevents out-of-memory errors on the executors and “straggler tasks”
Partitioners

- The partitioner defines **how records will be distributed** and thus which records will be completed by each task.

- Practically, a partitioner is actually an interface with two methods:
  - `numPartitions` that defines the number of partitions in the RDD after partitioning.
  - `getPartition` that defines a mapping from a key to the integer index of the partition where records with that key should be sent.
There are two implementations for the partitioner object provided by Spark

- **HashPartitioner**
  - Determines the index of the child partition based on the hash value of the key

- **RangePartitioner**
  - Assigns records whose keys are in the same range to a given partition
  - Required for sorting since it ensures that by sorting records within a given partition, the entire RDD will be sorted

- It is possible to define a custom partitioner

Partitioners and transformations

- Unless a transformation is known to only change the value part of the key/value pair in Spark
  - The resulting RDD will **not have a known** partitioner
    - Even if the partitioning has not changed
Using narrow transformations that preserve partitioning

- Some narrow transformations, such as mapValues, preserve the partitioning of an RDD if it exists.
- Common transformations like map and flatMap can change the key.
  - So even if your function does not change the key, the resulting RDD will not have a known partitioner.
  - Instead, if you don’t want to modify the keys, call the mapValues function (defined only on pair RDDs)
  - It keeps the keys, and therefore the partitioner, exactly the same.
  - The mapPartitions function will also preserve the partition if the preservesPartitioning flag is set to true.

The contents of this slide-set are based on the following references

- SQL Joins: https://www.w3schools.com/sql/sql_join.asp