

CS 455: INTRODUCTION TO DISTRIBUTED SYSTEMS

[THREADS]

The House of Heap and Stacks

Stacks clean up after themselves

But over deep recursions they fret

The cheerful heap has nary a care

Harboring memory leaks, hurtling to a crash

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Topics covered in this lecture

- Creation and Management
- Thread lifecycle
 - ▣ Creating and starting threads
- Stopping and interrupting threads
- Approaches to writing threads
 - ▣ Subclassing Threads vs Implementing Runnable



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THREADS AND ... HEAPS AND STACKS

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Threads and heaps

- For performance reasons, heaps may **internally subdivide** their space into per-thread regions
 - ▣ Threads can allocate objects at the same time *without interfering* with each other
 - ▣ By allocating objects used by the same thread from the same memory region?
 - Cache hit rates may improve
- Each subdivision of the heap has **thread-local variables**
 - ▣ Track parts of thread-local heap in use, those that are free, etc.
- New memory allocations (`malloc()` and `new()`) can take memory from **shared heap**, only if local heap is used up



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How big a stack?

[1/2]

- The size of the stack must be large enough to accommodate the **deepest nesting level** needed during the thread's *lifetime*
- Kernel threads
 - ▣ Kernel stacks are allocated in physical memory
 - ▣ The nesting depth for kernel threads tends to be small
 - ▣ E.g. 8KB default in Linux on an Intel x86
 - ▣ Buffers and data structures are allocated on the heap and never as procedure local variables



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How big a stack?

[2/2]

- User-level stacks are allocated in virtual memory
- To catch program errors
 - ▣ Most OS will trigger **error** if the program stack grows **too large too quickly**
 - ▣ Indication of an unbounded recursion
 - ▣ Google's GO will automatically grow the stack as needed ... this is very uncommon
 - ▣ POSIX for e.g. allows default stack size to be library dependent (e.g. larger on a desktop, smaller on a phone)
 - ▣ "Exceeding default stack limit is very easy to do, with the usual results"
 - ▣ Program termination



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THREAD LIFECYCLE

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Lifecycle of a thread

- Creation
- Starting
- Terminating
- Pausing, suspending, and resuming



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Thread: Methods that impact the thread's lifecycle

```
public class Thread implements Runnable {  
    public void start();  
    public void run();  
    public void stop();  
    public void resume();  
    public void suspend();  
    public static void sleep(long millis);  
    public boolean isAlive();  
    public void interrupt();  
    public boolean isInterrupted();  
    public static boolean interrupted();  
    public void join();  
}
```

} Deprecated, do not use



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Thread creation

- Threads are represented by instances of the Thread class
- When you extend the Thread class?
 - ▣ Your instances are also Threads
- We looked at the 4 constructor arguments in the Thread class



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Starting a thread

[1/2]

- Thread exists once it's been constructed
 - ▢ But it is *not executing* ... it's in a **waiting** state
- In the waiting state, other threads can *interact* with the existing **thread object**
 - ▢ Object state may be changed by other threads
 - Via method invocations



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Starting a thread

[2/2]

- When we're ready for a thread to begin executing code
 - ▢ Call the **start()** method
 - ▢ **start()** performs internal house-keeping and *then calls* the **run()** method
- When the **start()** method returns?
 - ▢ **Two threads** are executing in parallel
 - ① The original thread which just returned from calling **start()**
 - ② The newly started thread that is executing its **run()** method



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After a thread's `start()` method is called

- The new thread is said to be **alive**
- The `isAlive()` method tells you about the state
 - `true`: Thread has been started and *is executing* its `run()` method
 - `false`: Thread may *not be started* yet or may be *terminated*



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Terminating a thread

- Once started, a thread executes only one method: `run()`
- This `run()` may be complicated
 - May *execute forever*
 - Call *several other methods*
- Once the `run()` *finishes* executing, the thread has **completed** its execution



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Like all Java methods, `run()` finishes when it ...

- ① Executes a `return` statement
- ② Executes the last statement in its method body
- ③ When it *throws an exception*
 - ▣ Or fails to catch an exception thrown *to it*



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The only way to terminate a thread?

- ▣ Arrange for its `run()` method to **complete**
- ▣ But the documentation for the `Thread` class lists a `stop()` method?
 - ▣ This has a *race condition* (unsafe), and has been deprecated



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Some more about the `run()` method

- Cannot throw a **checked** exception
- But it can throw an **unchecked** exception
 - ▣ Exception that extends the `RuntimeException`
- A thread can be **stopped** by:
 - ① **Throwing** an unchecked exception in `run()`
 - ② **Failing to catch** an unchecked exception thrown by something that `run()` has called



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Pausing, suspending and resuming threads

- Some thread models support the concept of **thread suspension**
 - ▣ Thread is told to **pause** execution and then told to **resume** its execution
- Thread contains `suspend()` and `resume()`
 - ▣ Suffers from vulnerability to **race conditions**: **deprecated**
- Thread can **suspend its own execution** for a specified period
 - ▣ By calling the `sleep()` method



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But sleeping is not the same thing as thread suspension

- With true thread suspension
 - ▣ One thread can suspend (and later resume) *another thread*
- `sleep()` affects only the thread that executes it
 - ▣ Not possible to tell another thread to go to sleep



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But you can achieve the functionality of suspension and resumption

- Use `wait` and `notify` mechanisms
- Threads **must be coded** to use this technique
 - ▣ This is not a generic suspend/resume that is imposed by another thread



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Thread cleanup

- As long as some other active object holds a reference to the terminated thread object
 - ▣ Other threads can execute methods on the terminated thread ... retrieve information
- If the object representing the terminated thread goes *out of scope*?
 - ▣ The thread object is **garbage collected**



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Holding onto a thread reference allows us to determine if work was completed

- Done using the `join()` method
- The `join()` method
 - ▣ **Blocks** until the thread has completed
 - ▣ *Returns immediately* if
 - The thread has already completed its `run()` method
 - You can call `join()` any number of times
- Don't use `join()` to poll if the thread is still running
 - ▣ Use `isAlive()`



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STOPPING A THREAD

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Two approaches to stopping a thread

- Setting a flag
- Interrupting a thread



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Stopping a Thread: Setting a flag

- **Set some internal flag** to signal that the thread should stop
- Thread periodically **queries the flag** to determine if it should exit



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Stopping a Thread: Setting a flag

```
public class RandomGen extends Thread {  
    private volatile boolean done = false;  
  
    public void run() {  
        while (!done) {  
            ...  
        }  
    }  
  
    public void setDone() {  
        done = true;  
    }  
}
```

run() method investigates the state of the done variable on every loop. Returns when the done flag has been set.



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Interrupting a thread

- In the previous slide, there may be a *delay* in the `setDone()` being invoked & thread terminating
 - ▣ Some statements are executed after `setDone()` and before the value of `done` is checked
 - ▣ In the worst case, `setDone()` is called right after the `done` variable was checked
- **Delays** while waiting for a thread to terminate are *inevitable*
 - ▣ But it would be good if they could be minimized



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Interrupting a thread

- When we arrange for thread to terminate, we:
 - ▣ Want it to *complete its blocking method* immediately
 - ▣ Don't wish to wait for the data (or ...) because the thread will exit
- Use `interrupt()` method of the `Thread` class to **interrupt** any *blocking method*



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Effects of the interrupt method

- Causes blocked method to **throw** an **InterruptedException**
 - ▣ `sleep()`, `wait()`, `join()`, and methods to read I/O
- Sets a **flag** inside the thread object to indicate that the thread has been interrupted
 - ▣ Queried using `isInterrupted()`
 - Returns true if it was interrupted, even though it was not blocked



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Stopping a thread: Using interrupts

```
public class RandomGen extends Thread {  
    public void run() {  
        while (!isInterrupted()) {  
            ...  
        }  
    }  
}
```

`radomGeneratorThread.interrupt()`



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The Runnable interface

- Allows **separation** of the *implementation* of the task *from the thread* used to run task

```
public interface Runnable {  
    public void run();  
}
```



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Creation of a thread using the Runnable interface

- Construct the thread
 - ▣ Pass runnable object to the thread's constructor
- Start the thread
 - ▣ Instead of starting the runnable object



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Creation of a thread using the Runnable interface

```
public class RandomGenerator implements Runnable {  
    public void run() { ... }  
}  
  
...  
generator = new RandomGenerator();  
Thread createdThread = new Thread(generator);  
createdThread.start();
```



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When to use Runnable and Thread

- If you would like your class to inherit behavior from the Thread class
 - ▣ **Extend** Thread
- If your class needs to inherit from other classes
 - ▣ **Implement** Runnable



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If you extend the Thread class?

- You **inherit behavior** and **methods** of the Thread class
 - ▣ The `interrupt()` method is part of the Thread class
 - ▣ You can `interrupt()` **if you extend**



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Advantages of using the Runnable interface

- Java provides several classes that handle threading **for** you
 - ▣ Implement pooling, scheduling, or timing
 - ▣ These require the **Runnable** interface



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But what if I still can't decide?

- Do a UML model of your application
- The object hierarchy tells you what you need:
 - ▣ If your task needs to subclass another class?
 - Use Runnable
 - ▣ If you need to use methods of Thread within your class?
 - Use Thread



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Threads and Objects

- Instance of the Thread class is just an **object**
 - ▣ Can be passed to other methods
 - ▣ If a thread has a reference to another thread
 - It can invoke *any method* of that thread's object
- The Thread object is not the thread itself
 - ▣ It is the set of methods and data that *encapsulate* information about the thread



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But what does this mean?

- You cannot look at the object source and know *which thread is*:
 - ▣ Executing its methods or examining its data
- You may wonder about which thread is running the code, but ...
 - ▣ There may be many possibilities



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Determining the current thread

- Code within a thread object might want to see that code is being executed either:
 - ▣ By thread represented by the object or
 - ▣ By a completely different thread
- Retrieve reference to current thread
 - ▣ `Thread.currentThread()`
 - ▣ Static method



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Checking which thread is executing the code

```
public class MyThread extends Thread {  
  
    public void run() {  
        if (Thread.currentThread() != this) {  
            throw new IllegalStateException  
                ("Run method called by incorrect thread ...");  
        } /* end if */  
  
        ... Main logic  
    }  
}
```



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Allowing a Runnable object to see if it has been interrupted

```
public class MyRunnable implements Runnable {  
  
    public void run() {  
        if (!Thread.currentThread().isInterrupted() ) {  
            ... Main logic  
        }  
    }  
}
```



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The contents of this slide-set are based on the following references

- *Java Threads*. Scott Oaks and Henry Wong. . 3rd Edition. O'Reilly Press. ISBN: 0-596-00782-5/978-0-596-00782-9. [Chapters 3, 4]
- *Operating Systems Principles and Practice*. Thomas Anderson and Michael Dahlin. 2nd Edition. ISBN: 978-0985673529. [Chapter 4]



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