Plan for Today

Recursive descent or predictive parsing
- example predictive parser
- FIRST and FOLLOW sets revisited
- constructing a predictive parser table

Syntax-directed translation

Example Predictive Parser

```c
void S() { switch(tok) {
  case NUM:   Mesh(); eat(EOF); break;
  default:    error();
}

void Mesh() { switch(tok) {
  case NUM:
    num_nodes = NUM.val; eat(NUM);
    NodeList();
    num_elem = NUM.val; eat(NUM); break;
  default:    error();
}

void NodeList() { switch(tok) {
  case NUM:   break;
  case NODE:  Node(); NodeList(); break;
  default:    error();
}

S -> Mesh EOF
Mesh -> num NodeList num ElemList
NodeList -> e | Node NodeList
Node -> node num real real  // node_id, x, y
ElemList -> e | Elem ElemList
Elem -> tri num num num num // elem_id, 3 node ids
Elem -> sqr num num num num // elem_id, 4 node ids
```

FIRST and FOLLOW sets

nullable(X)
- X is a nonterminal
- nullable(X) is true if X can derive the empty string

FIRST
- FIRST(ε) = {ε}, where ε is a terminal
- FIRST(X) = \bigcup FIRST(rhs)_i, where X is a nonterminal and X -> rhs_i
  - union all of FIRST(sym) on rhs up to and including first nonnullable

FOLLOW(Y), only relevant when Y is a nonterminal
- look for Y in rhs of rules (lhs -> rhs) and union all FIRST sets for symbols after Y up to and including first nonnullable
- if all symbols after Y are nullable then also union in FOLLOW(lhs)

Constructing the Predictive Parser Table

Algorithm
for each X -> gamma
  for each T in FIRST(gamma)
    table[X,T] = X->gamma
  if gamma is nullable
    for each T in FOLLOW(X)
      table[X,T] = X->gamma

S -> Mesh EOF
Mesh -> num NodeList num ElemList
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Syntax-directed translation

One-pass versus multi-pass compiler
- What do we need for MiniJava?

In a predictive, or top-down, parser
- do actions in functions for nonterminals
- example: storing off the number of nodes and elements in the mesh grammar parser

In a shift-reduce, or bottom-up, parser
- each reduction leads to an action being performed
- example: SableCC builds a parse tree using actions associated with each grammar rule
- syntax-directed translation is equivalent to performing an action on internal nodes in the parse tree during an in post-order traversal
- example: interpreter you wrote for PA1